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Towards Secure and Privacy Preserving Collision Avoidance System in 5G Fog based Internet & Vehicles

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Abstract

Current avoidance systems mainly focus surrounding entities including the pedestrians, 'he cyclists are assumed to use a different avoidance system for their aff y. Vehicle speed is reported as one of the major factors that causes ... h sev re road accidents that affect other entities on the road. In response, several solutions have been implemented to control the causalities of over spee 'ing ranging from speed camera, speed detectors to car avoidance systems. However, those solutions have not significantly improved the rate of truffic acci ents and their impact. Additionally, the current solutions do not ensure "in ely notification of all the road users (surrounding vehicle drivers ped strights or others) that can alleviate crash causalities in case of fatal ' affic . rivents. The fifth generation (5G) cellular network is predicted to cher me the current limitations of Internet-of-Vehicles (IoV) by offering fast i w latency and reliable connections to enable IoV based applications. Fc, cor puting has also been proposed to complement IoV by bringing computational tities in nearby proximity of the vehicles. 5G based fog vehicul. " n' wor's is a new paradigm that empower real-time and low latency sees for telligent Transportation System (ITS). In this paper we proposed ¿ secure a d privacy-preserving collision avoidance system in 5G fog based IoV. The 2 devices are used to collect speed violation report (TVR) sent by the

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vehicles' speed sensors. The fog nodes aggregate multiple TV (s, v m, the signatures on the TVRs and broadcast anonymous notifications to other entities in the vicinity. The protocol makes use of certificateless aggregal signeryption coupled with pseudonymous technique as the building blocks to consure authentication, integrity, confidentiality and privacy preservat; respectively. The batch verification technique is used by the fog devices to allow simultaneous TVRs signature verification for a timely response. The authorization of reporting speed sensors is both guaranteed by the to ption-based information along with the digital signature to discard all the bog is TVRs. The analysis of the protocol confirms its lightweightness and "ticiency.

Keywords: 5G cellular networks; for computing; intelligent transportation system; security; collision avoidance sys an.

1. Introduction

The recent report of World Health Organization released in May 2017 records about 1.24 million peor e who d e each year as a result of road traffic crashes. Road traffic injuries are cite ' o be the leading cause of death among young people, aged betw on 1, to '9 years and could be the seventh major cause of death by 2030 [1,2]. The report ranges vehicle speed as of one of the main factors which ause and accidents along with the drink driving, non usage of belt and distraction. Approximately 90 % of the road accidents occur in low-and middle-in. $m\epsilon$ countries due to inadequate road safety infrastructure as well as poor t^{*} affic mana, ement system. In most high-income countries, around 20 % of 10 all the 'r fic e cidents are caused by exceeding the speed limit [3][4]. However, ir these high income countries, most of the roads are equipped by cameras ϵ hd speed detectors in order to monitor and latter on identify the drivers who violated the permitted speed limit. Though those cameras have significantly impr ved the traffic congestion issues, those infrastructures have not achieved inficant result for preventing or alleviating the causalities caused by the traffic accidents [5][6]. Furthermore, traffic accidents caused by over-speeding vehicles

have more causalities whereby other entities in the vicinity such as pe_{-}^{*} ortrians, motorcyclists, or other vehicles might be involved in the acciac +. Therefore,

²⁰ systems for warning over-speeding vehicles would be criti al to the improvement of road safety by alleviating the causalities of traffic accelents of used by overspeeding vehicles.

Currently, fixed speed detectors are found on the roads α developed countries, their main role is to capture any vehicle that exc. As the speed limit on the

- ²⁵ particular point and probably fine the over-speeting vehicle/driver. Conversely, those fixed infrastructures seem not to have any offect on the over-speeding issue since the drivers can use the navigat. In systems or warning signs to know where those cameras are located and reduce then speed for the sake of not being fined [7]. Nevertheless, using the vehicular communications, the entities in the
- ³⁰ vicinity of an over-speeding vehic, on bo alerted or warned to avoid major causalities in case of traffic ac idents.

Lately, intelligent transportation systems (ITS) has received the attention of both the industry and a significant through various projects [8]. The main goal is to offer a variety of the bad server as through the cloud based vehicle to vehicle

- ³⁵ (V2V) and vehicle t infrast. ture (V2I) communications. Actually, V2V is set to be more us ul f r loc lized emergency services while V2I is considered for non critical ervices [?]. However, the cloud based vehicular networks' solutions preser num, "ous issues related to transmission of significant real-time traffic data ro, the roads infrastructures to the cloud servers which cause time
- ⁴⁰ delays and we y costly in terms of bandwidth [10]. Additionally, the IEEE 802.11, Long-Term Evolution (LTE) standards, that were initially proposed for vehiciding communications, revealed scalability and mobility support issues for vancular communications [11]. Thus, the 5G cellular networks is predicted to empower TS based services through its features including massive bandwidth, massive connectivity and reduced latency [12] [13].

F cently, a new computing paradigm referred as fog computing was presented. This computing architecture stretches the conventional cloud computing and respective services to the network level. This paradigm offers several

features including low latency, extensive geo-distribution, p sitio . a preness, enhanced mobility coupled with real time service processes [14]. Contrary to the convectional central cloud-based systems, the fog b used model allows the sensors to transmit the data to nearest fog devices. Thus, fog devices can perform computation on the collected data and help for necision making [15]. While the integration of fog computing and 5G cellular networks of ne to fruition, pri-

- ⁵⁵ vacy and security issues should be carefully address d. This appeals for an innovative design of secure and privacy preserving protocol for potential critical services in 5G fog based vehicular network. In this paper, we present a secure and privacy preserving protocol it. collision avoidance system with the feature of fog devices that will enable the data seconded by the vehicles speed
- sensors to be aggregated and sent to by neighboring entities in the vicinity of the violating vehicle. This would reduce the causalities of traffic accidents caused by over-seeding vehicle. To the best of our knowledge, this is the first study that address specific secure, and privacy preserving issues for collision avoidance through an over preding reporting system in 5G fog based vehicular
- networks. For a such oplicatio to gain people's consideration and motivate the stakeholders for 'mpleme.' .tion, security requirements should be satisfied. Beyond the routine security objectives such as confidentiality and authentication; it is crucia' to prote ' the real identities of some entities in the system that send/receive the reporting messages including the vehicles, the pedestrians or
- The motore ch. s. For instance, the transmitted messages reporting the overspeed veh. 'a nould not be accessed by unauthenticated entities on the roads. The system should satisfy mutual authentication between the speed sensors, the fog at 'a s, the roadside clouds as well as the trusted entities. Also, the system should be feasible by demonstrating an acceptable lightweightness in terms of communication and computation overhead.

Taking into account the heterogeneous architecture of 5G cellular networks, the μ comising merits of fog computing, the security objectives to be achieved, we are encouraged to design a secure and privacy preserving protocol that enable collision avoidance by reporting the over-speeding vehicle to address the

- ⁸⁰ aforementioned challenges. The motivation behind this pape⁻ are ¹⁰ ws:
 - Current vehicles are already equipped with several ensors but the information collected by the sensors are not fully utilized in real time. Depending on the type of the sensor, the collected information represents different levels of sensitivity. For example, over speeding *P* and subled braking data could be used differently. In this research, *we confider the speed record*ing sensor that contains information which could be collected, analyzed, and transmitted to other entities (venibles, procestrian, motorcyclist) in the same vicinity as the speeding vehicle formations.
 - The privacy and security properties of the speeding vehicle, the roadside clouds, the fog devices, even other relicles in the same range and direction should be met.

Consequently, the contribut. γ or uns paper are threefold:

- We first present an FAPA ation model for secure and privacy preserving collision avoidance costem in G fog based Internet of Vehicles which allows the vehicles' steed senses at the send the recorded traffic violation reports (TVR) to fix d frig no less. The fog nodes aggregate the received reports and anony mously no lifty the entities in the vicinity. TVRs are then sent to other fog no des and to road side clouds. We define the attack model of our application model and the requirements to be met by the proposed protocial in terms of security and performance.
- We esign a secure and privacy preserving collision avoidance protocol in 5G. * ased Internet of Vehicles based on the techniques of certificateless agg. gate signcryption, pseudonymous and batch verification techniques.
- we provide the analysis of the proposed protocol in terms of security and performance. We further evaluate the performance of our protocol through computational delay, transmission overhead and simulation.

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The remainder of the paper is organized as follows. We first p_1 cont the related work, the system model and cryptographic primitives r_1 constructing the proposed scheme in section 2, 3 and 4 respectively. We prepare the design of the proposed protocol in section 5. We discuss security and reformance of the proposed protocol in section 6 and draw the conclusion in section 7.

2. Related Work

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In this section, we first overview the convertion, 'V'.NETs to the 5G enabled Internet of Vehicles. Then we outline the fogenetworking architecture and then provide available literature for secure collis. In avoidance systems in vehicular networks.

2.1. VANETs, 5G-Enabled Internet . even cles

Vehicular ad hoc netwoks 'VANE.'s) extends the conventional Mobile ad hoc networks (MANETs) [16]. In 'ANETs, the key entities represent the vehicles, the fixed infrastruc' as on the roads named road side units (RSU) and

- hicles, the fixed infrastruct are on the roads named road side units (RSU) and a third party named T usted Ai thority (TA) responsible for the registration, certification and revolution o. ", e entities participating in the VANETs architecture. Conventional VA".ETs offers two major communication techniques using the dedicated short rang. communication (DSRC) which are; vehicle to vehicle
- (V2V) and v nicle ~ infrastructure (V2I) communications [17] [18]. Several application were projected to be implemented through the VANETs framework, how ere the computational overhead of the applications in VANETs necessite es sufficient computation capabilities that appealed for the mixture of VANETs vith cloud computing [19] [20]. The feasibility of VANETs integrated
- ¹³⁰ to cloud computing also called *Internet of Vehicles* was adopted by several resistances archers 21]. Nevertheless, some drawbacks were drawn among the available technologies for vehicular networks. For instance, the IEEE 802.11p has been proved by the researchers to suffer from mobility support [22].

Furthermore, the Long Term Evolution (LTE) in the 4G cellular networks does not provide required latency required for the vehicular networks [23] [24] [25]. Hence, the key features of 5G cellular network in terms ϵ_{\perp} late α_{ν} massive connectivity, spectral efficiency and data rate establish a prom. ing paradigm for the (ITS) [26].

- 2.2. Fog Computing
- Lately, several researchers have documented the s de of a t of fog computing and presented the IoT applications which could benc⁺ from the fog networking including smart city, smart health care and c hart grid [27] [28]. The merits of fog computing which aims at relieving the computation and communication burden on the cloud computing, provid. - an intermediate layer between the cloud and mobile/fix devices to offer smooth and law latency delivery services,
- were adopted [29] [30]. Yi et al., introdu ed possible latency sensitive application areas including real time video set a res, content delivery and caching, and big data analysis [31]. The authors pointed out two main issues related to resource management and computation on a reding. Dantu et al., gave a comparison of fog
 computing and the conventional cloud computing based on energy consumption

and latency [32].

In [33] Tao et al. press *e , the integration of fog networking and cloud computing to avai' 5G- nabl d Vehicle to- Grid (V2G) networks which would facilitate several V2G ... *v ork. The authors also showed that fog networking ¹⁵⁵ can achieve 26 % a. * 90% for time response reduction for users and data traffic respectively ... *v vehicular communications, fog computing has been adopted by several rese rchers as a promising technology to implement real time services [34] [3^r] [36]. In [37], Lingling et al. presented a secure and privacy preserving

- navig 'io' scheme, he fog devices use traffic information collected from the vehicles to compute optimal route. Thus the vehicles can get optimal routes continuously
- from the fog devices. For edge computing in internet of things including internet of vehicles, several research have been proposed recently. Min Chen and the device the device the task offloading issue in ultra-dense network in order to minimize the delay while gaining the battery life [38]. In [39], the authors

suggested a joint mobility-aware caching and Small Base Stat on (' \downarrow ,) density placement scheme based on the user mobility. For the IoV, a μ_{ν} , architecture named Cognitive Internet of Vehicles (CIoV) that focus on intra vehicle, intervehicle and beyond-vehicle network was proposed [40]. \geq ang et al., proposed

- software defined network (SDN) based concept the cender of traffic safety by detecting the driver's fatigue detection [41]. In [42], when it al., proposed an Edge-Cognitive-Computing architecture for smart-he. "theore system that help to analyze and monitor the health of patients councilities, computing. Though the above articles tackle several issues related to edg. computing, there is no compre-
- ¹⁷⁵ hensive, secure and privacy preserving sc. ome in the literature that addressed the security and privacy threats for objective avoidance through over-speeding scenario in 5G fog based vehicular networks.

2.3. Collision Avoidance System

Collision avoidance systems (i.e. bedestrian or other entities on the road) can be divided into three groups: Infrastructure based systems that focus on availing innovative transport infrast actures that allow the separation of cyclists, pedestrians or vehicles. Pace we collision avoidance system aiming at reducing the damages after a collision such as road bumpers. Active collision avoidance system that use detection systems and sensors to alert drivers of potential accidents [43] [4] [4] [4] The active collision systems can also be divided in three groups: windows, radar and vision based technologies. The radar and vision based systems are affected by numerous limitations such as recognition latency, weather, line of light and so on. Thus, wireless based technologies can overcome the logitations of radar and vision based technologies. Recently, several manu-

- fa current static ced to equip the vehicles with a set of sensors including the vehicle s beed sensor (VSS), carbon monoxide sensor (CMS), alcohol sensor (ALS), Gas leakage sensors (GLS), vehicle noise sensors (VNS), ect,... [46]. The data collecter through the sensors are still strictly used within the vehicles and do not in efft neither the surrounding entities (vehicles, pedestrian, motorcyclist,...)
 - nor the transportation authorities which would use it for several purposes [47].

Jang et al., presented a fixed sensor based intersection collisien warm system [48]. Their protocol makes use of vehicle's location, speed and time collected from several sensors located at the intersection to warm for an evolution to the intersection. However, their article focused on collaboration in the intersection.

- collision warning system rather than over-speed r' portion. Additionally, the authors did not address the security and privacy concerns of the proposed protocol. For traffic violation monitoring, Mallissety et 1 [49], proposed a traffic violation monitoring system in VANETs. Howe, r. their protocol was not built under the heterogeneous 5G fog based architect. And the security and privacy features were not deeply investigated.
- ²⁰⁵ features were not deeply investigated.

Considering the above discussion (*) is obvious that the proposed solutions in the literature do not address the rep. (*) ng of speed violators that can relieve the damages in case of accidents, (*) also the proposed solutions do not take advantage of 5G fog based free mowork, which provide low latency services for

210 ITS. As a result, it is arguable that "here is not direct research in the literature that designed a secure ap", "ivacy preserving collision avoidance system in an heterogeneous 5G fog b sed Inte net of Vehicles.

3. System Mod 's a' d D sign Goals

This section includes the system model, the communication model, the adversary model, the security requirements and finally the design goals.

3.1. Systen. ' odel

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T¹ is system mode is made by a master overviewer TA, the road side cloud (RSCs), $\arg r$ gator fog device (AFD) and the speed sensors (SS) that are incory orated. The vehicles as shown in Fig. 1. We describe the role of each entity in the following:

• Transportation Authority (TA): TA is a fully trusted public agency that registers all entities in the system (SS, AFD and RSC) and provides cryptographic materials during the system initialization.

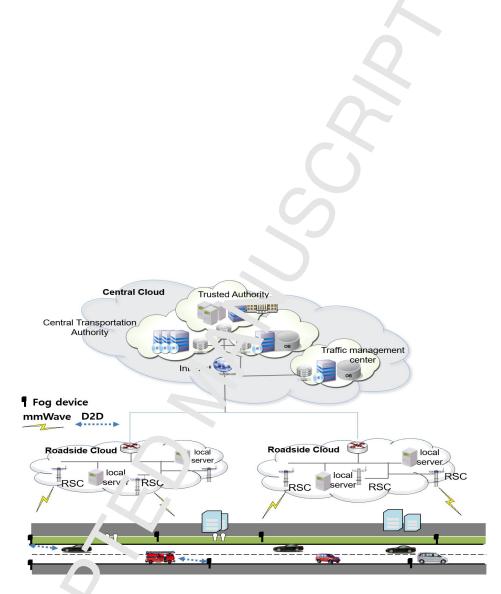


Figure 1: System Architecture

• Speed Sensing Nodes (SS): SS also denoted as $SS = \{SS_1 \cup S, \dots, SS_n\}$ (where *n* symbolize the number of sensors that send the coefficient violation report on a given time) are in-built speed sensor in the chicle. These sensors use available GPS data to know the vehicle speed limit within a location. These sensors could also be portable device such smart phones or tablets. The sensing nodes are accountable for sensing the traffic violation reports to the aggregator fog devices (Ai \Im s).

- Aggregator Fog Devices(AFD): Like a " "htwe.," server, AFDs are devices fixed all along the roads with computing etc. age and communication capabilities. For example, they can be a red on the road light poles. The AFDs are connected wirelessly to spece an asing nodes. The main tasks for the AFDs are; to collect the "VRs, regregate them and perform signature verification on aggregated TV.'s. The over-speed warning messages are broadcast to all authorize ' sume inding entities, then forwarded to other AFDs and eventually to the RSCs.
- RoadSide Cloud RSC): h 3Cs are databases fixed along the roads and communicate v .th the ^F Ds. The RSCs store the traffic violation report (TVR) sent b other AFF s. Then the TVR can be sent to the TA for further legal purst .t. The PC C can also if needed broadcast the valid (aggregated and verimed) TVRs. RSCs are assumed to be connected to an electricity power 5C erator with sufficient computational capability.
- Vehicle. We assume that all the vehicles are equipped with speed sensors and onnected to GPS in order to know the speed limit in a given area.
 The OP J collects the data through the SS and sends them to AFD using D2. Yor mmWave communications. All vehicles are supposed to register with the TA where periodical inspection usually takes place. Beyond the conventional identifier of the vehicle including the Electronic License Plate (ELP) or Electronic Chassis Number (ECN), each vehicle in the system is assumed to have a 5G unique identifier (5GID), similar to the

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mobile phone subscriber identification module (SIM) r imber within the conventional 3G and 4G cellular networks.

255 3.2. Communication Model

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Motivated by the 5G cellular networks architecture, the proposed 5G fog based Internet of Vehicles is made by the following portponents:

- Heterogeneous Networks: This network aims a "chieving high data rate and network capacity for the 5G fog bas. ⁴ network. Therefore, two alternatives may help to get the mentioned producties through smaller cells which increase the spectral efficience [150]; and using the mmWave spectrum would offer high data ration since to operates within the range of 30–300 GHz and 1–10 mm for the spectrum and wavelength respectively [51].
- D2D Communications: Line contraction would enable the speed sensors to communicate with each tog device within the licensed cellular bandwidth without constatering the Base stations. In the 5G fog based vehicular networks, the contraction between the vehicles, speed sensors and fog devices can be done through D2D communication or mmWave technology.
- In the protocol's s_y ter model, the first phase concerns the communication between the S^c (s_y red sensor) nodes and the AFDs (aggregator fog device). The communication between those two entities is made possible through the D2D or mmWr re teranologies. However, we also assume that the existing and inexpensive WiFi to behave would be used in remote areas where the 5G cellular

275 networks are not available. On the other hand, the second phase in our system module contract as the communication between the AFDs and the RSCs. Since the constance to tween those entities could be significant, the communication channel count the a wired or wireless link that offer low delay along with high bandwidth. In our system model, we adopt the following assumptions:

• We assume that the TA has sufficient storage capabilities, strongly protected and hard to be compromised by an adversary. TA is responsible

for generating all key pairs both for the SSs, the AFDs and ∞ during the system setup phase. TA can also maintain the list c. compromised entities (vehicles, fog nodes or RSCs).

- There exist one RSC in an area of 600 meters of adiu. Luch is trustworthy. RSC is responsible for keeping the TVI s for fur her purposes. We assume that there are at least four AFDs between the RSCs.
 - Every SS sensor node communicates w.' exact' one or several AFD. AFD is responsible for the aggregation . 'd verification of TVRs. AFD broacast TVRs and send them to . 'Os for further purposes.
 - We assume that all the entitie in a system have clocks for generate time stamps and to check time value by of exchanged messages. One of the existing solution is the use of GPS artellite for time source synchronisation [52].

²⁹⁵ 3.3. Adversary model

The full adoption of which application partly relies on how the security threats are handled; thus it is important to study and address all the means which the adversary can use in confusion the whole system. Within our system model, we assumed that the RCSs and AFDs are honest, but curious entities. Nevertheless, there could be an adoption provide the AFDs that could eavesdrop on the TVRs. Additionally, an adversary A could also access the personal information of the vehicles the wight the AFDs databases. The adversary could also be able to launch different attacks such as false injection attack to threaten the integrity of the Wirks. The unauthorized access to the vehicle's personal information voluble if to privacy violation. Impersonation and masquerading attacks would a no lead is traffic jam since other vehicles in the same vicinity as the violating vehicle would be taking preventive actions to avoid any damages which could be crused by the violator.

3.4. Design Goals

3.4.1. Security Objectives

- Privacy preservation: The SSs involved in sensing the TVRs should be protected from revealing their sensitive person, data such as their respective identities. Additionally the real nontitie of other entities in the system that process TVRs should be preserved.
- *Mutual authentication*: The ADFs and the SSs should authenticate each other to avoid that an external and multicious user would interfere and jeopardize the system.
- Data Confidentiality and Integria: All transmitted TVRs should be delivered unaltered.
 - Authorization: The TVn, should be sent by legitimate SSs only and processed by legitima. AFDs
 - Key Escrow Lesil Ace: The trusted authority and the motor department which generate the Veys should not have the full private keys of all the entities in the ovstem. Therefore, even though those key generation centers are compromised, the adversary can not get the full private keys of the entities.
 - *l'rac ıbility*: The TA should be able to reveal the real identities of all the p_a icip cing entities in case of dispute.

4.2. Pe formance Objectives

• communication and verification Overhead: The secure protocol should be efficient in terms of communication overhead and offer suitable processing latency. A significant number of TVRs should be verified, aggregated in a very short interval.

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In this section, we describe the design goals of the proposed p. *ocol which comprise the security objectives and the performance of jectives.

- Robustness: Though some of the entities may be intrud d, the array sent from the SSs to the ADFs should not be accessed.
- Lightweight: Vehicle speed sensors and fog devices have mited power and storage capabilities. Thus, the proposed <u>potocolon</u> nould have low computational cost.

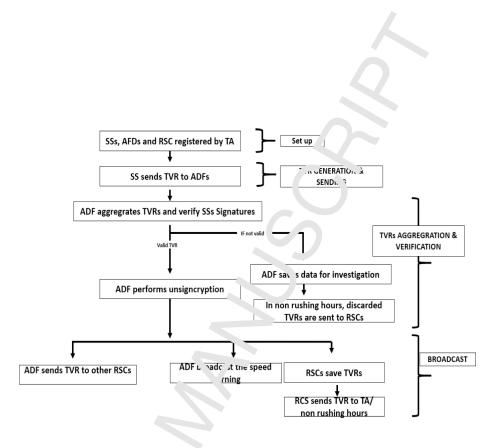
3.5. Overall Protocol Description

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From the above described architecture, our source and privacy preserving protocol for collision avoidance system is made ν . the following sub phases as shown in Fig. 2:

Initialization: TA sets up its n. st r secret key and its corresponding public key. Each vehicle provides its real identity and TA generates the corresponding pseudo identity for the SS from which a partial signing key is generated. AFD and RC, provide their real identities and TA assign the partial private keys. All the entities in the system including the SSs, AFDs, RSCs register with the TA.

- Traffic violation report generation and sending: When a vehicle enters a particular is of we assume that the speed sensor registers the specific accepted beed limit. Whenever a vehicle goes beyond the specified speed limit, the OBC best the information from the speed sensor, composes message call which the OBU signerypts. The message is sent to the closest AFD.
- FVT s ag regation and verification: Upon receiving the TVRs, the AFD perto. s the TVRs aggregation and verification. Since a single AFD can receive multiple TVRs simultaneously, AFD aggregate the TVRs for fast precessing. Later on the AFD performs the signature verification on aggregated TVRs. This will help to discard all the bogus TVRs which might have been sent.



re 2: System Flowchart

• Traffic violation Rep. * Jeaconing: In case everything holds, the AFD reformulates ne TVRs by removing personal information in the message that can compromise the security and identity of the sender and broadcast the speet. We ming. Simultaneously, the AFD forward the message to the closest if TD which will also broadcast the reformulated TVRs. In case the foccing AFD is close to an RSC, it will also forward the message to the security and regional traffic authority (TA). The light is described protocol preserves the privacy of the road violators, the integristic for further pursuit.

4. **Preliminaries**

In this section, we described the certificateless scheme of signcryption (CLSC) [03] and bilinear paring [54] which are considered as our building blocks. In our

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- ³⁷⁵ construction we have used pseudonymous identity for the vehicle $s_{\rm r}$ and sensors to preserve their real identities. The CLSC scheme was ${\rm ado}_{\rm r}$ and to suit our system model for the following reasons:
 - 1. CLSC does not use certificate for authorization 1. _____ne system can avoid computational overheads caused by ce tific te revocation, storage and distribution.
 - 2. In CLSC, the full private key of the users are no generated by the TA. Since the system can be deployed country. 'do through regional transportation authorities, this would prevent the pllapse of the whole system in case one regional transportation au.' ority is compromised.
- 385 3. CLSC performs both signature and e application in a single step, this helps the protocol to be lightweig' which is a crucial feature for the adoption of such application.

4.1. Bilinear Maps

Let \mathbb{G}_1 and \mathbb{G}_2 be two eyens properties of some large prime order q. The bilinear ³⁹⁰ map $\hat{e}: \mathbb{G}_1 \times \mathbb{G}_1 \to \mathbb{G}_2$ so 'isfies t' e following properties:

- Bilinear: $\hat{\mathbf{e}}(a^{\mathsf{T}}, bQ^{\mathsf{T}} = \hat{\mathbf{e}}(P, Q)^{ab}$, for all $P, Q \in \mathbb{G}_1$ and all $a, b \in \mathbb{Z}_q^*$.
- Non-dege erate: It . is a generator of G₁ then ê(P, P) is a generator of G₂.
- Cor out a le: There is an efficient algorithm to compute $\widehat{\mathbf{e}}(P,Q)$ for any $P\ Q\in {\mathbb{G}}_*$

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4.2. Ce. fica cless Aggregate Signcryption

In the following section, we describe the main functions of the proposed precord ased on the CLSC in [53]. Those functions include the setup, partial public key generation, partial private key generation, full private key generation, sign ryption, aggregation, aggregate-verification and aggregate-unsigncryption.

4.2.1. Set up [Cas.Setup()]

Let \mathbb{G}_1 be a cyclic additive group with a prime order q on alliptal murve, and P be an arbitrary generator of \mathbb{G}_1 . Let \mathbb{G}_2 represents a cyclic multiplicative group satisfying a bilinear map where $\hat{e} : \mathbb{G}_1 \times \mathbb{G}_1 \to \mathbb{G}_2$. Cas. So imply a secured by the TA and output the parameters as follows:

1. Randomly selects a master private key $s \in \mathbb{Z}_{+}^{s}$ ar ¹ compute the master public key $P_{Key} = sP$. Note that the master privale key s is kept securely by the overwiewer TA.

2. Chooses four hash functions $H_1 : \{0 \ 1\}^* \quad \mathbb{Z}_q^*, H_2 : \{0,1\}^* \rightarrow \{0,1\}^n$

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with *n* being the bit-length for the plan. Txts to be secured. $H_3: \{0,1\}^* \to \mathbb{G}_1$ and $H_4: \mathbb{Z}_q^* \to \mathbb{G}_1$.

3. Set $Cas.params = (\mathbb{G}_1, \mathbb{G}_2, P, q, I_{ub}, H_1, H_2, H_3, H_4)$

Cas.PUK() is computed by the for VID_i to generate a partial public key as follows

- 1. VID_i randomly c. Set $x \in \mathbb{Z}_q^*$ as a secret value and generates a partial public key $K_i = x P$.
- 2. VID_i forwards '; ide tity and the partial public key (VID_i, K_{ib}) to TA.

4.2.3. Partia Prive's Key Generation [Cas.PVR()]

- $_{420}$ Cas.PI \checkmark () is run by the TA to generate a partial private key as follows:
 - TA che ses y_i ∈ Z^{*}_q and compute additional partial public key for VID_i is K_a = y_iP. Then the complete public key for the vehicle becomes (n_i K_i).
 - 2. TA generates the partial private key $D_i = y_i + s * PID_i$ where $PID_i = H_1(/ID_i)$. (PID_i, D_i) is sent to VID_i in a secure manner.
 - 4.2.4 Full Key Set Algorithm [Cas.Skey()]

Cas.Skey() is performed by the user VID_i after the verification of the partial private key provided by the TA:

- 1. VID_i checks the legitimacy of the partial private key by verifine whether $D_iP = K_{ia} + P_{key}H_1(VID_i)$ and set the full private key (x_i, D_i) .
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- 4.2.5. SignCryption Algorithm [Cas.SignE()]

Given a message m_i , an aggregation keyword \triangle , *sas.SianE()* is executed by VID_i along with the receiver ID ID_R . VID_i perference the ollowing:

- 1. VID_i selects $r \in \mathbb{Z}_q^*$ and generates $T_i = rP$.
- 435 2. Compute $Z_b = rY_{rb}$.
 - 3. Compute $Z_a = r(Y_{ra} + P_{key}PID_i)$.
 - 4. Compute $h_a = H_2(ID_R||Y_{ra}||Y_{rb}|| \triangle_{||} ||Z_b||Z_a)$
 - 5. Compute $F_i = h_a \bigoplus m_i$.
 - 6. Compute $h_b = H_3(ID_R||Y_{rc} |V_{rb}| \land |T_i||F_i||PID_i||K_{ib}||K_{ia}).$
- 440 7. Compute $h_c = H_4(\triangle)$.
 - 8. Compute $\alpha_i = D_i h_c + r h_b r_i h_c$.
 - 9. Return the ciphertext $C_i = (T_i, F_i, \alpha_i)$.
 - 4.2.6. Aggregation Algo, "hm [C s.Aggr()]

Taking the receiver $I^{\prime} ID_R$, Cas.Aggr() is executed by the receiver through

- 445 the following steps:
 - 1. Generat's $\alpha = \sum_{i=1}^{n} \alpha_i$
 - 2. Retur $= (T_1...T_n, F_1...F_n, \alpha).$
 - 4.2.7. Aggrey 'on-Verification Algorithm [Cas.AggrV()]

T le re eiver ID_R runs Cas.AggrV() by computing the following:

- 450 1. Compute $h_b = (ID_R ||Y_{ra}||Y_{rb}|| \triangle ||T_i||F_i||PID_i||K_{ib}||K_{ia})$ for i = 1, ..., n
 - 2. Con pute $h_c = H_4(\triangle)$
 - 3. Compute the verification by running

$$\hat{e}(\alpha, P) = \hat{e}(\sum_{i=1}^{n} K_{ia} + P_{key}PID_i, h_c)\hat{e}(\sum_{i=1}^{n} T_ih_b)\hat{e}(\sum_{i=1}^{n} K_{ib}, h_c).$$

The correctness is verified as follows

 $(\alpha, P) = \hat{e}\left(\sum_{i=1}^{n} \alpha_{i}, P\right)$ $= \left(\sum_{i=1}^{n} (D_{i}h_{c} + rh_{b} + x_{i}h_{c}), P\right)$ $= \left(\sum_{i=1}^{n} D_{i}h_{c}, P\right)\hat{e}\left(\sum_{i=1}^{n} rP, h_{b}\right)\hat{e}\left(\sum_{i=1}^{n} x_{i}P, h_{c}\right)$ $= \left(\sum_{i=1}^{n} D_{i}P, h_{c}\right)\hat{e}\left(\sum_{i=1}^{n} T_{i}, h_{b}\right)\hat{e}\left(\sum_{i=1}^{n} K_{i}b, h_{c}\right)$

 $= \left(\sum_{i=1}^{n} K_{ia} + P_{pub}PID_{i}, h_{c}\right) \hat{e}\left(\sum_{i=1}^{n} T_{i}h_{b}\right) \hat{e}\left(\sum_{i=1}^{n} {}^{V_{c}}_{ib}, h_{c}\right) \text{ If all the equations}$

⁴⁶⁰ hold, it outputs true, otherwise false.

4.2.8. Aggregation-Unsigncrypt Algorithm [Cas. A. rV()]

In case Cas.AggrV() holds, the receiver \mathcal{P}_R performs the following to unsigncrypt the ciphertext:

- 1. Compute $Z'_b = x_r T_i$
- 2. Compute $Z_{a}^{'} = D_{r}T_{i}$

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- 3. Compute $h'_a = H_2(ID_R||Y_{ra}||V_{rb}|| \Delta ||T_i||Z'_b||Z'_a)$
- 4. Compute $K_i \bigoplus h'_a$
- 5. Finally, it output $\{m_i\}_{i=1}^n$ The correctness is verified as follows

 $\mathbf{m}_{i}^{'}=F_{i}\oplus h^{'}$

 $= \mathrm{H}_2(PID_i||K_{ia|} \lor b|| \land ||T_i||Z_b||Z_a) \oplus m_i \oplus h_a'$

- $=\mathrm{h}_{a} \in m_{i} \in h_{a}^{'}$
- $= r_{\iota}$

5. Prop. v. Protocol

In this section we present a secure and privacy preserving collision avoidance system for 5G fog based Internet of Vehicles which is made by following main sub protocols: system initialization, traffic (speed) violation report(TVR) generation and sending, TVR aggregation and verification and TVR broadcast. The list of notations within the protocol are found in Tab. 1

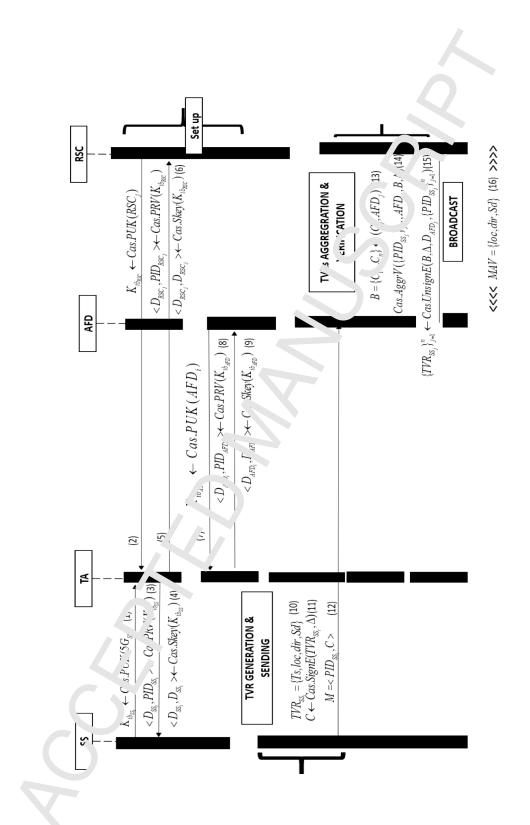




	Table 1: Notations and Descriptions
Notation	Description
$5G_{SS_i}$	Unique 5G identity for each sen or i
TA	Trusted Authority
s_{TA}	Trusted Authority private ey
RSC_j	Roadside cloud's server
AFD_i	Identity of aggregator fog dev. ?
\bigtriangleup	Location based keywora
TVR	Traffic(speed) violation
K_{ib_j}	Partial public key \sum an entity j
D_i	Private key (
s_i	Full private key \cdot entity i
PID_i	Pseudo ide. h_{J} for an entity i
Ts	time (
Sd	vehicle speed
dir	v.m., direction
loc	whicle 1 cation
TVR	Trafac (speed) violation report
\mathbb{G}_1	Elliv cic curve group with the same order q
$P \in \mathbb{G}_1$	A generator of \mathbb{G}_1
VAM	Violation alarm message

5.1. Initializ ion

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W this this retup phase, TA generates the general parameters and the other entities register to TA. Note that some steps such as the generation of the TA's master servet key are not described in Fig. 3:

• *step 0.* TA selects an elliptic curve group \mathbb{G}_1 of order q and a generator $P \in \mathbb{G}_1$. TA computes the master secret s_{TA} and public key P_{TA} by running Cas.Setup() and set $\langle s_{TA}, Cas.params \rangle \leftarrow Cas.Setup()$, then publishes the parameters Cas.params.

- Step 1. Each speed sensor SS_i runs $K_{ib_{SS_i}} \leftarrow Cas.PU \ (5C_{SS_i})$ to generate its partial public key and sends it to TA.
- Step 2. RSC_j generates $K_{ib_{RSC_j}} \leftarrow Cas.PUK(R C_j)$ as a spartial key and sends it to TA.
- Step 3. TA generates $\langle D_{SS_i}, PID_{SS_i} \rangle \leftarrow as.P' R(K_{ib_{SS_i}})$ as the partial private key of SS_i and sends it to SS_i .
- Step 4. SS_i set its full private key by running $\sim S_i, D_{SS_i} \leftarrow Cas.Skey(D_{SS_i}) >$.
- Step 5. TA generates $\langle D_{RSC_j}, PI_{\kappa SC_j} \rangle \leftarrow Cas.PVR(K_{ib_{RSC_j}})$ as the partial private key of RSC_j and sends in γRSC_j .
- Step 6. RSC_j runs $\langle s_{RSC_j}, D_{RSC_j} \leftarrow Cas.Skey(D_{RSC_j}) \rangle$ to set its full private key.
- Step 7. Each aggregator fog ¹ vice AFD_i generates $K_{ib_{AFD_i}} \leftarrow Cas.PUK(AFD_i)$ as its partial public key and sends it to TA.
- Step 8. TA gener. $cos < D_A \circ_{D_i}, PID_{ADF_i} > \leftarrow Cas.PRV(K_{ib_{AFD_i}})$ as the partial private key of A. J_i and sends it to AFD_i .
 - Step 9. $A \cap D_i$, $ns \in s_{AFD_i}, D_{AFD_i} \leftarrow Cas.Skey(D_{AFD_i}) > \text{to set its}$ final prive key.

The TA de em nes also the formats of TVRs sent by the SSs.

505 5.2. T' iffic vii +ion report generation and sending

A the disc sor SS_i will use the GPS information to know the acceptable sr \ldots limit \perp a given area. For instance, in a school zone the speed limit is 1 prmally 0 km/h. In such environment any vehicle driving beyond that speed 'limit \ldots a cause SS_i to generate a TVR and forward it to AFD_i . Suppose a speed sensor SS_i of a vehicle v_i records a TVR, it performs the following:

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- Step 13. SS_i generates $\langle C \leftarrow Cas.SignE(T \ (R \ _{S_i}) \rangle)$
 - Step 14. SS_i sends $M = \langle PID_{SS}, C \rangle$ to Ar.

5.3. Traffic violation report aggregation and orification

The aggregation protocol is responsible for the aggregate multiple ciphertext Ms into a single M. This is very important to cause some highways can have ⁵²⁰ up to 16 lanes, 8 lanes in one side an ¹ 8 lanes in the opposite side. Even in the cities, we can easily find an 8 \cdot ¹2 lates road with 4 lanes or 6 lanes each side. Thus, hundreds of vehicles can violate the speed limit in few milliseconds. The system should be able to ag_8 regate those TVRs and report them on time. This is very useful for timely speed violation reporting. Suppose for a set of

- ⁵²⁵ $TVR_{SS_i} = \{T, loc, dir, \ s\}$ gen ``ated by $n < SS_1.....SS_n >$, we can achieve the aggregation of several ``4 `* $Aggr_M = < PID_{SS_1}, ...PID_{SS_n}, C_1, ...C_n >$. ADF_j performs the foll wing to aggregate the TVRs.
 - Step 13. FD_j takes $C_j = \{C\}_{j=1}^n$ and outputs the aggregated ciphertext $B = \{C_1, \dots, C_{n_j} (C_j, RSC_j).$
- Ste₁ 14. FD_j runs $Cas.AggrV(\{PID_{SS_j}\}_{j=1}^n, RSC_j, B, \Delta)$ for signature b.tch ven ration.

If the sign ture verification holds,

• Step 15. AFD_j outputs $\{TVR_{SS_j}\}_{j=1}^n \leftarrow Cas.UnsignE(B, \triangle, D_{AFD_j}, \{PID_{SS_j}\}_{j=1}^n)$ \therefore recover the TVRs.

535 5.4. Traffic violation report broadcast

Step 16. After the RSC_j has successfully recover the traffic fieldion reports TVR_{SSi} = {T, loc, dir, Sd}, it reformulates the messa e and broadcast the violation alert message VAM = {loc, dir Sd_f = bill an only contains the location of the vehicle, the direction and spectrum s we noted in the assumption, the vehicles on board units have display facilities and able to approximately show the position of the violating vehicle on the screen. Note that the VAMs can be forwarde 1 to RSC_j can play the role of AFD_j if the fielding vehicle is closer to RSC_j than AFD_j.

545 6. Security and Performance Analy is

In this section, we evaluate the proposed protocol in terms of security goals, computational cost and commune tion cost.

6.1. Security Analysis

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We discuss in this sec. In the security goals for overspeed reporting in 5G fog based vehicular networks as set in section 3.4.1

- Privacy p servation The proposed protocol guaranties the identity protection of the ontities participating in the system. First, an adversary can rot optice the identity of the SS_i or the AFD_j through eavesdrop₁ inc because there is no plain text within the transmitted TVR form SS_i to AFD_j since SS_i sends M =< PID_{SS}, C > with < C ← Ca .Sigr E(TVR_{SSi}) >. Also, during the registration phase, the SS_i or AFD_j is provided a partial private key with a pseudo identity PID_i = H₁('S_i). Given that the partial private key D_i = y_i + s * PID_i with r1D_i = H₁(VID_i), which contains the pseudo identity is sent securely to the requesting entity, the adversary cannot reveal the real identity of the SS_i by eavesdropping. Additionally, assume the fog devices are compromised by the adversary, the adversary would only get the SS_i pseudo

identity which can not reveal it real identity. Therefor, we on "m that the proposed protocol achieves identity preservation.

- 2. Authentication: The authentication between SS_i a' d AFL upon sending 565 a TVR is guaranteed by the signcryption on the TV?. Aft ϵ , generating a speed violation message $TVR_{SS_i} = \{Ts, loc, (ir, Sc^i), SS_i \text{ performs sign-}\}$ cryption on the message as $\langle C \leftarrow Cas.Sign_{\perp}(TVI_{SS_i}) \rangle$. Only the entity with valid full private key can unsigner, t the TVRs. Note that the adversary can not have the full private key of a speed sensor SS_i be-570 cause even the TA does not have the full $_{\rm F}$ ivate key of the SS_i . TA only generates the partial private key or the entities by running Cas.PVR(). Thus, entity authentication is revided by the certificateless aggregate signcryption technique. The secur ', of the signature depends on the unforgeability of CLCS scheme i. i. " an prively chosen message attacks [53]. 575 Consequently, we endors that the designed protocol guaranties entity authentication.
 - 3. Authorization: In t¹ proposed protocol, an unauthorized SS_i can not send any TVRs. I just the p otocol prevents the malicious users outside the RSC zone to g nerate a "VR. As described in section 5.2, SS_i composes a traffic viol just in less ge TVR_{SS_i} = {Ts, loc, dir, Sd} where Ts, loc, dir and Sd represent the sime stamp, the location, the direction and the speed respectiblely. S^C will then generate a zone secret value △ corresponding to each 'loca' on based on the RSC sites. Thus, this will prevent an adversary outs. 'P in RSC zone to generate TVRs. Additionally, even though an interpret in the zone generates a TVR, the TVR signature verification is not registered by the TA. Therefore, the proposed protocol achieve entity authorization.

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4. Confidentiality and Integrity: The speed sensor SS_i generates a TVR and signcrypts it as $\langle C \leftarrow Cas.SignE(TVR_{SS_i}) \rangle$. Note that within the signcrypt function Cas.SignE(), the ciphertext contains $C_i = (T_i, F_i, \alpha_i)$ where T_i and F_i accomplishes the functionalities of message encryption

and α_i the duties of digital signature. In the propored proce γ , only the legitimates AFDs can perform the unsigneryption of γ through the computation of T_i , F_i and α_i . Therefore, since the pertification less signature is proven to be secure under adaptively chosen ciphe text (U dD-CC2) [55], we confirm that the proposed protocol achieves represent the proposed confidentiality and integrity.

- 5. Key Escrow Resilience: The massive connective, of the 5G cellular networks required distributed systems to avoid key management burdens. Thus, the applications under the 5G fogored architecture ought to satisfy the Key Escrow resiliency. In the proposed protocol, the speed sensors generate their partial public ' we by computing Cas.PUK(). The SSs sends the partial public key to TA of the full private key of SS is computed by SS_i after the verification of the partial private key generated by the TA. Therefore, we confirm that the proposed protocol achieves key escrow resilience property.
- 6. Traceability: In the proposed protocol, the TA generates the entity pseudo identities PID = H₁(S⁻) and saves the hash values in a table. Therefore, in case of dis_F at the ⁷ A is able to reveal the real identities of the entities in the dispute by checking the corresponding hash value of the pseudo identity which the response of the proposed protocol guarantie, the traceability of the participating entities.

6.2. P rforman. Analysis

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In bi section, we provide the performance analysis of the proposed protocol bosed on the computational and communication cost.

6.z. Computational Cost

1 CLSC [53], three main operations are executed; the scalar multiplication overlated in group \mathbb{G}_1 , the exponentiation operation that is calculated in the

group \mathbb{G}_2 , and lastly the pairing operation. Those three operation a. respectively denoted as T_{mul} , T_{exp} and T_{pair} . However, the proposed construction only performs T_{mul} and T_{pair} . To measure the computation cost of the proposed protocol, we made use of an MNT curve along with the Taul pairing $\hat{e} : \mathbb{G}_1 \times \mathbb{G}_2$ on the curve, the embedding degree is 6 and the i_i is regimented by 160 bit [56]. The implementation was done on a desktop computer with 3.5GHz, core i-5, 16GB RAM using the pairing based library in a. ¹ the Miracl library. The execution time are depicted in Tab. 2.

Table 2: Measurement of		
Notation	Operations	time (ms)
T_{pair}	Bilinear pai. 'ng	4.5
T_{mul}	Point se 'ar mutiplication	0.6
T_{exp}	Modular e. 2011, ntiation	1.4

- In the proposed secure and privacy preserving collision avoidance protocol, whenever a SS_i senses ε speed violation, it executes 6 T_{mul} to sign a TVR as described in section 4.2. On the other hand, the receiver AFD_j performs 4 T_{pair} to aggregate, lerify and unsignerypt the TVRs as shown in section 4.2.7. Fig. 4 shows the tool 1 cost of signing one or multiple TVRs. As shown in the figure, the time for verifying multiple TVRs is stable due to batch verification
- technique which is used for TVR signature verification. Thus, the proposed protocol vould or instance require 18 ms to verify 1000 TVRs. For the signing process we as time that a SS_i signs one TVRs at a go which cost 3.6 ms.

Table 3: Computational cost of proposed protocol				
Operation	$\mathrm{Cost/ms}$			
$6T_{mul}$	3.6ms			
$4T_{pair}$	18ms			
	21.6			
	Operation $6T_{mul}$			

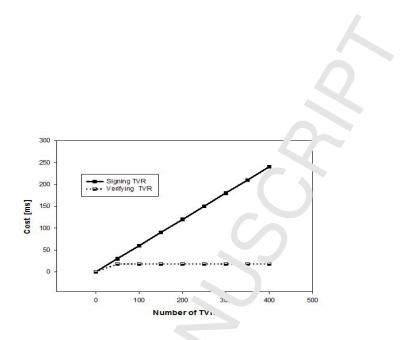


Figure 4: Overall Sign/Verica. " Cost of TVR

6.2.2. Communication Cost

- In this section, we provide the an lysh of communication cost for the secure and privacy preserving collision avoid here system. The communication cost includes the cost of transmitting the TVRs from the SS_i to the AFD_j . We first emphasize on the transmistion overhead caused by the signeryption which was performed on the 1.7°. In f_{11} , the sizes of the elements are $64 \times 2 = 128$ bytes and $20 \times 2 = 40$ bytes for \mathbb{G}_1 . The sizes for the hash functions is 20 bytes and the other element such as the time stamps have a 4 bytes size [57]. In the proposed prote 1 the signerypt function Cas.SignE() contains the ciphertext $C_i = (T_i, F_i, \alpha_i)$ where $T_i = 20$ bytes, $F_i = 60$ bytes and the signature $\alpha_i = 56$ bytes. The size of the a raw warning message is 40 bytes according to the society
- of automotive ingineers. Therefore the size of raw TVR is 40 bytes where as a secure TV , size is 136 bytes [58].

ℓ 2.3. Sⁱmutation

In this subsection, simulation experiments are provided in two folds. First we investigate the impact of the TVR size on the computation offloading. Secondly, we iv vestigate the impact of vehicle's speed, vehicle's density and fog device's density on the loss ratio.

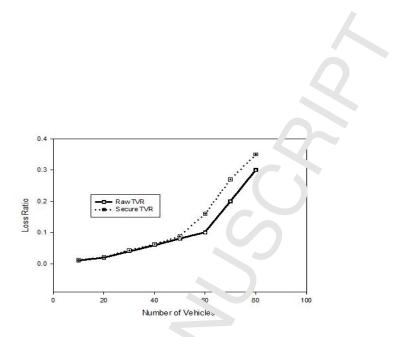


Figure 5: Impact of Number of vehicle " the average loss ratio

6.2.3.1 Impact of TVR Size on Comput tional Offloading

Assume that each AFD needs a comp. *ation amount M_i to execute a task T_{task} of a TVR of size TVR_{size} . The computing resource within edge/fog computing

- ⁶⁶⁰ is set to 25 GHz and the computing resource for vehicles OBU is set to 10 GHz.As described in subsection 3.1, the vehicle's OBU could also be assumed as a mobilephone fixed in the web cles. The distance between one RSC to another is 600 m in which we set 3 bace stations (BS). We adopt *Energy Optimal (GSI)* as a benchmark Ligories we'recause it helps the user (OBU) to connect with the
- best channel r ga. "less of the delay performance. Energy Optimal (GSI) is the standard 3C^T LTE protocol for handover [59]. The size of the secure and non secure T^{*} Rs *e* described in Tab. 4 as calculated in subsection 6.2.2. The details of the sn. "llation setting are set as described and suggested in [60].

V 2 evoluate the size of the TVR with and without security features on 670 cc², utatic. 1 offloading. As shown in Fig. 6, the larger the TVR size, the 1 onger the task would take. However the difference is not considerable looking at unover the difference is not considerable looking at unover the effects which could be obtained for non-secure TVR. For instance, for a 60 sytes TVR message, the task execution is is 0.5 milliseconds for unsecured message and 0.59 milliseconds for secured VTR. The overall increase is of 9 % which is not considerable. In this subsection, we neglected the investigation

that output the impact of TVR size on energy cost. This is because we consider that a vehicle's OBU is not subject to energy issues since it keep. On recharging whenever the vehicle's engine is on. Additionally we also assume that the fix fog devices can also be connected to an electricity generation resource.

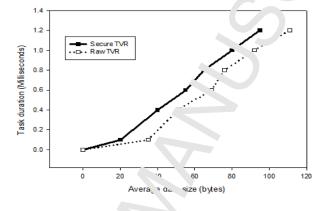


Figure 6: Impact o. "VR size on the task duration

680 6.2.3.2 Impact of vehic's's dens ty and speed on loss ratio

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In addition, we make us of VANET-SIM simulator to enable vehicle mobility; then for network limu, 'ior we used ns-3 simulator [61]. We later set our system scenario using ... IEEE 802.11p platforms for the 5G cellular network which is predicted 'o range from 1 Gbps for highly populated roads to a maximal transmission range of 9-10 Gbps [62].

The ns-3 L set using the Friis equation that describes the propagation of signal as $I_r = \frac{P_t G_t G_r \lambda^2}{(4\pi l)^2}$ with P_t being the transmission power, then G_t and G_r represent the antenna gains, λ represents length of the wave and l represents the distance from the receiver to the transmitter [63].

"" ermore, we downloaded a map from OpenStreepMap website [64]. In our imulation, each vehicles is randomly released and can move randomly wit¹ in the map. The speed of the vehicles is set from 10 to 40 m/s that is "7 to 144 km/hr. . Tab. 4 describes more details concerning the settings for

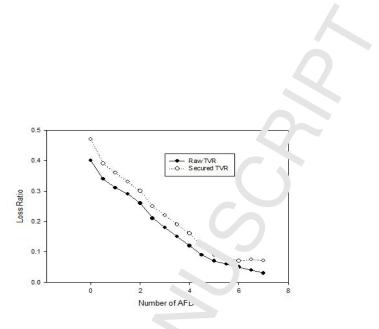


Figure 7: Impact of ADF on the verage loss ratio

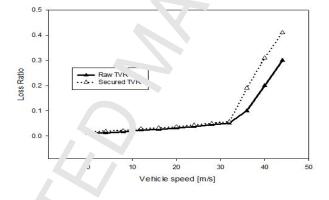


Figure . "mpact of vehicle speed on the average loss ratio

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We later call lated the message loss ratio (AV_R) as follows [65]:

$$AV_{R} = \frac{\sum_{i=1}^{DE} \sum_{j=1}^{TVR} \sum_{k=1}^{FD} (T_{F} - T_{C})}{\sum_{i=1}^{DE} TVR} + Aver$$

v here $DF_{i}TVR, FD, T_{F}, T_{R}$ and Aver represent the density of the vehicles, the number of TVR sent by v_{i} , the number of AFD_{j} within the simulation area, the t me when v_{i} forwards a TVR message to AFD_{j} , the time AFD_{j} receives a $\cdot \cdot R$ message from v_{i} and the average aggregation verification time that AFD_{j} authenticates/aggregates and verify the TVRs respectively.

Table 4: Simulation settings				
Tools/Parameter	Value/Specificatio.			
Mobility generation tool	VANETS M 2.02			
Network Simulation tool	ns-3			
Trans range	1GE)s			
Number - of - vehicle	15			
Simulation time	200 mm			
Wireless protocol	80. ¹¹ 1p			
TVR propagation interval	6 . ^			
Departure interval	$20 \ sec$			
RSC radius	600 m			
Number of AFD within 2 1 $^\circ$ C	3			
mobility model	$shortest \ path$			
Message size for TTTh	40 bytes			
Message size for secure. TVR	136 bytes			

Fig. 7 shows the a grage loss ratio according to the number of available AFD. We can see that as the aumber of fog devices increases, the average loss ratio decreases. The gef re, the fog based computing within the cellular networks overcome the communication overheads of convectional vehicular networks. In
Fig. 5 and 8, we show the average loss ratio based on the number of vehicles and the vehicle speed. As we can see, the raw TVR which has a smaller size performs both if that a secured TVR. For instance, in Fig. 8, the loss ratio starts increasing significantly when the violating vehicles go beyond a speed of 34 m/s which is 122 km/h. In Fig. 5, the performance of the proposed protocol for secured TVR when the number of vehicles is less than

7 Thus in urban areas where the speed limit range from 60 to 80 km/h, the computational and communication overhead caused by the security features do not ; fect significantly the performance of the overall system.

7. Conclusion

- ⁷¹⁵ While we wait the fulfillment of 5G fog based internet of vehicles, security and privacy should be carefully addressed for ITS applications. Therefore, we proposed in this paper a secure and privacy preserving collision avoidance system in 5G fog based internet of vehicles. The fog dences are used to collect speed violation report (TVR) sent by the vehicle speed sensor. The batch
- verification techniques allow the fog devices \uparrow verify 1 ultiple TVRs simultaneously. The features of certificateless aggreg, \uparrow signaryption scheme that offer both the encryption and digital signature in a line step were adopted to securely transmit the speed violation reports. The proposed protocol is suitable for distributed systems since it meets 1. to the neutrine security goals as such
- ⁷²⁵ authentication, confidentiality and ...+egri, : but also the key escrow resilience . The performance analysis in terms of ...>mputation and communication overhead confirms its efficiency.

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SI on Future Generation Computer Systems

Title: Towards Secure and Privacy Preserving Collision Avoidance System in 5G Fog based Internet of Vehicles

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SI on Future Generation Computer Systems

Title: Towards Secure and Privacy Preserving Collision AvoiCance System in 5G Fog based Internet of Vehicles

Highlights

- Proposed a secure and privacy-preserving model for collision avoidance system in 5G fog based Internet-of-Vehicles (Uv)
- Constructed a secure protocol using signerypt.on, pseudonymous and batch verification techniques
- Security and performance analysis with respect to various metrics