# **Accepted Manuscript**

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PII: DOI: Reference:	S0167-739X(18)30670-8 https://doi.org/10.1016/j.future.2019.01.022 FUTURE 4715
To appear in:	Future Generation Computer Systems
Received date :	28 March 2018
Revised date :	27 November 2018
Accepted date :	13 January 2019

Please cite this article as: U. Noor, Z. Anwar, A.W. Malik et al., A machine learning framework for investigating data breaches based on semantic analysis of adversary's attack patterns in threat intelligence repositories, *Future Generation Computer Systems* (2019), https://doi.org/10.1016/j.future.2019.01.022

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# A Machine Learning Framework for Investigating Data Breaches Based on Semantic Analysis of Adversa. v's Attack Patterns in Threat Intelligence Repositories

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### Abstract

With the ever increasing cases of cyber data breaches, the manual process of sifting through tons of security logs to investig te cyper-attacks is error-prone and timeconsuming. Signature-based deep search plutices only give accurate results if the threat artifacts are precisely provided. With the purgeoning variety of sophisticated cyber threats having common attack patterns and utilizing the same attack tools, a timely investigation is nearly impossible. The is a need to automate the threat analysis process by mapping adversary's Tactics, Techniques and Procedures (TTPs) to attack goals and detection mechanisms. In this poper, a novel machine learning based framework is proposed that identifies cyber t' reats be sed on observed attack patterns. The framework semantically relates threats and " "Ps ' stracted from well-known threat sources with associated detection mechanis is to form a semantic network. This network is then used to determine threat occurrent s b for hing probabilistic relationships between threats and TTPs. The framework i train, <sup>4</sup> using a TTP taxonomy dataset and the performance is evaluated with three reported in threat reports. The framework efficiently identifies attacks with 92% at uracy and low false positives even in the case of lost and spurious TTPs. The average detection time of a data breach incident is 0.15 seconds for a network trained vith 133 TTPs from 45 threat families.

Keywords: Cyl er Threa' Intelligence, Data Breach Investigation, Tactics Techniques and Procedure . In licators of Compromise, Belief network, Latent Semantic Indexing

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Preprint submitted to Elsevier

November 23, 2018

#### 1. Introduction

The pervasiveness of high-speed Internet connectivity has attracted . In ge majority of businesses to move their sensitive and confidential information and upmactions to the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. This transition has made business operations flexible and end for the clouds. The billion data records were breached in a total of 1792 incident in the first half of the year 2017, over 1.9 billion records were breached in 918 incidents [2 including the massive data breach of Equifax where 148 million consumers had their performance is use common attack patterns to compromise their target. Due to this reacon, he security community is paying more attention towards sharing and consuming Cyber for the first line the provide the provide the security defense against data built busines.

The massive size of CTIR and the constant or laught of new Advanced Persistent attack's behavioral signature. Currently, there are  $ap_{\mu}$  eximately 1 billion threat indicators publicly available on Hail-a-TAXII [4]. Sin. arry, IBM X-Force [5] reports thousands of malware on a weekly basis. The Verizon's Data Leach Investigations Report (DBIR) [6] reports millions of data breaches including such credit card credentials. Along with these reports, a network administrator has loc,  $^{1}$  y available log files [7] containing useful threat data, e.g., DNS [8], packet captur, [2], en ail [10] and IP logs [11]. Among these different types of reports, Structured Thr. at Information Expression (STIX) [12] encoded CTIRs are more comprehensive news they record the details of cyber breaches as attack observables, indicators, TTPs, n. idents, threat actors, campaigns, exploit targets and Course of Actions (COAs). To investigate data breach incidents, so far, a major focus of the security cor munit, has been on sharing and identifying indicators and observables. These are low-lovel the at artifacts comprising of IP addresses, domain names and file hashes. Unfortunated they have a very short lifespan with respect to threat defense as they are susc otible to change having fewer chances of being reused. The attacker constantly changes IF addresses by buying new attack servers and domain names. While cyber da' a breach incidents are caused by such malware and APTs that share attack patterns 's 1. Ps. The TTPs represent the attacker's actions for breaching an organization's net ork. As the indicators and observables change rapidly, similarly the goals and targe s of attackers also change frequently. Ironically the attackers' TTPs remain the same and e re-used over and over again with little innovation. This problem was observed in the data breach incidents of some notable organizations like Sony [13]. Target [14] ar He me Depot [15] where millions of customers were affected. In these incidents, the adv. "se les' employed the same TTPs to breach into the network. The greatest ar ount of pain can be inflicted upon the attacker by identifying TTPs in the network as 1 therely achieving a timely investigation.

In this reason work, we propose a novel and automated data breach investigation frame ork that exploits adversary's attack patterns, i.e., TTPs in CTIR. The framework general es a Tareat-TTP-Detection (TTD) network by semantically correlating threat incidents and themselves to identify the existence of a TTP in a network. The TTD set, and twork is further augmented to investigate the most probable threat incident famine based on detected TTPs in a network, by adopting a machine learning based



Figure 1: Pyramid of Pai Mou. 11

probablistic analysis. The proposed framework pliably detects data breach incidents with an accuracy of 92% and low false portions even in the case of lost and spurious TTPs.

The organization of the paper follow this equence: The problem statement is discussed in the second section. Based on the problem statement we describe our research methodology and contributions in section three. Section four discusses the related work. The proposed framework design is detailed in the fifth section. The working of this framework is explained in the sixth section with the help of a financial RAM scrapping malware family data breach case on the effectiveness and efficiency of the framework are evaluated in the seventh section. Finally, in the eighth section, the research work is concluded along with future perspectives.

#### 2. Problem Statemen<sup>+</sup>

Cyber data breach inclusts are caused by such malware and APTs that share attack patterns as IOCs. There are six levels of IOCs as defined in the Pyramid of Pain model [16, 17] shown in figure 1. The IOCs present in the first three levels are atomic indicators comprising file has. APS, domain names, network and host artifacts. Low-level IOCs, also termed *tech aical* the pattern (IDS) and spam filters on email servers. In the upper part of the pyram. The reare more generalized threat indicators related to the behavioral attack signed ures of threats such as exploit toolkits, malware and TTPs. The high-level IOCs, also termed as *tactical* threat intelligence is consumed by incident responders for investiga ion. Lardening defenses by upgrading systems and policies. Interestingly the indicators in the lowest part of the pyramid can be easily identified and extracted from the threat reports due to their fixed format. Therefore, they can also be easily defended by polying intrusion detection and firewall rules. Currently, the emphasis of the security community is to employ low-level IOCs for sharing and investigating a data breach incident. Unfortunately, these low-level IOCs are only useful for a short time or



Figure 2: Credential Compromise TTPs for Three Threat Groups

for immediate use. The attacker constantly changes 'P addresses by buying new attack servers and domain names. On the other han' high lovel IOCs are for a long-term use and provide sophisticated defense against cyber usreats. Thus by identifying high-level IOCs, i.e., TTPs the attacker is forced to c' and bis attack patterns.

To illustrate the significance of TTPs in the investigation process, in figure 2 an example of three famous threat groups, i.e.  $\neg G-1, 14$  [18], TG-3390 [19] and TG-4127 [20] is provided. All these groups compromise fictors' credentials but using different attack TTPs. The threat group TG-1314 for the victim's remote access tools by compromising credentials of the network's endpoint management platform. The threat group TG-3390 installs a keylogger and employs a publicly available credential dumper tool to get password hashes. The threat group TG-4127 compromises credentials via spear phishing. A well-known example of the relate 1 to TG-4127 is the Democratic National Committee's Gmail-based email  $\varepsilon$  courts by compare also different. The goal of the threat group TG-3390 steals industriation in the victim's remote access facility while the threat group TG-3390 steals industriation of the explorate and property. The threat group TG-4127 targets government and military networks for espionage and cyber warfare. This shows that by identifying the presence of a particular threat group related TTP in a network provides a distinct advantage to the analyst in explicitly identifying the cyber threat family.

Our problem suffer lent is thus as follows. Current techniques to diagnose a data breach incident use low level IOCs which are not useful due to their short life span. On the other *k* and if high-level IOCs, i.e., TTPs are used to investigate a data breach incident, they called a more accurate detection. Therefore our problem is to identify TTPs based on the low level threat artifacts observed in a network with the help of appropriate machine learning algorithms.

### 3. Research Methodology and Contributions

To access the problem statement presented in the previous section, here we describe ou. result is methodology from which we will derive our solution and accomplish our research challenges. To accurately investigate cyber threats using adversary's TTPs, two main challenges need to be addressed. The *first* challenge is that to  $\therefore$  ntify TTPs of an attack, the security analyst has to perform certain detection 'lech nisms that are frequently not specified in the CTIR. The *second* challenge is that  $\neg \neg$  's are rarely referenced using standard identifiers in CTIR. They are mostly reported is unstructured, human understandable textual descriptions that make it difficulted in correlate attack incidents of same threat group based on similar TTPs due to symptomy s and polysemous words. Manual searching for correlated TTPs is a tedious, time is not prone process which is nearly impossible due to the massive size of CT1R.

The *first* step of the methodology addresses the challenge of ident 'ying the existence of TTPs. The detection mechanisms of our framework are on tainer, from the publicly available (ATT&CK) taxonomy provided by MITRE [1] that documents adversarial Tactics, Techniques and Common Knowledge. The firm worl generates a semantic network of Threats, TTPs and Detection menchani.  $\gamma s$  (T $\gamma \tau$ ) by correlating threat incidents among themselves to identify the existence of a TP in a network. The second step of the methodology addresses the challenge C refere long TTPs in CTIR. The TTPs in the CTIR are semantically mapped into up T1.2 using their appropriate labels in the ATT&CK taxonomy using Latent Semantic Inouring (LSI). LSI, not only groups documents using semantically similar words but. 'so groups semantically similar words in a document to identify themes or topics of the documents. The TTD is further augmented to diagnose the most probable threat family, bay su on detected TTPs in a network, by adopting a probablistic machine learning base analysis using belief networks between threats and TTPs. The advantage of this approach is that it outperforms sophisticated classification methods, treats all predictor at vibutes independently and is useful for very large datasets with missing data [22].

To evaluate the working of the proposal framework, we constructed a benchmark dataset from the ground truth  $d\gamma'$  (threat, TTP and detection mechanisms) available on ATT&CK taxonomy [21]. Threat a tifacts were compiled based on the threat incidents reported by multiple sources, i.e., IBM X-Force [5], Symantec [23], FireEye [24] and CrowdStrike [25]. The entities were populated into the prototype system and automated investigation were populated to determine the presence of attacks.

Our research has the 'ollowing contributions:

- 1. This work is the usu f its kind that presents a case for emphasizing the analytics of high-level adversarial uTPs. Previous work has primarily focused on the identification of low level atomic indicators that are trivial for the attacker to change.
- 2. The results provided highlight the important relationship between security incidents, tacket and a tifacts in a way that machines can identify these connections with cert in r cobe bilities. Henceforth this research paves the way for cyber security investigation. with partial or incomplete information.
- 3. A we 1-know a standard dataset of unstructured cyber threat incidents has been thore ighly stilled and dissected to construct a comprehensive vocabulary of structured in TP... The existing threat taxonomies are either too specific to a particular comain, e.g., web attacks or are too verbose and descriptive to be useful for machine 1 arning
- 4. Our research motivates the development and usage of a common vocabulary for T. 1.5 in CTIR. The SIRS module allows security analysts to identify common TPs in different documents even if they have different texts but have the same

meaning. At the same time, this also gives SIRS the ability to ... over TTPs mentioned in one document but not the other. We illustrate this in the evaluation section by showing how the proposed system identifies missing TTF. I three documents (ATT&CK, IBM X-Force, Symantec) that concern the same "breat (AXIOM Group). These TTPs can be derived or essentially imported inclusion one lataset into the other thereby automatically "augmenting" the latter. This is hence an important contribution to our work.

5. The results prove that the cyber threat classes and their associated incidents utilize a finite and typically deterministic collection of TTPs, which characterize them. If this set of TTPs can be systematically documented then "be threat groups can be automatically profiled and ultimately stopped. Machines can help investigate these "telltale" threat group signatures. This is not surprising as threat groups constitute real people who specialize in particular tactics and techniques which evolve over time.

#### 4. Related Work

Cyber threat diagnosis based on learning  $_{\rm P}$  terns from CTIR is a relatively new domain. We did not find any directly related res. rch work that can be compared to the proposed scheme. However, our research occuefits from several existing research domains. First, we studied the different state if the art solutions that deal with cyber threats. Second, we studied research were where CTIR have been used to forecast cyber incidents, suggest investment in data security and operational risk management in financial sectors. Third, we studied how the proposed sources and enriching intrusion detection knowledge between to detect attacks and vulnerabilities. Fourth, we studied the most state-of-the-rest solutions for malware, APT and intrusion detection using machine learning technique

A machine learning bas d cyber inreat detection model to identify the seven top security threats in cloud om atin, and their remediation is proposed in [26]. The presence and type of att k is <sup>1</sup>et mined by training the machine learning model with a limited set of artifact 'hat depict the activity pattern of the cloud. The artifacts are performance logs of C.<sup>2</sup>U, st. "age media, network usage patterns of the hypervisor and the guest operating  $\gamma_{\rm e}$  em. In [27], the authors propose an abstraction layer over the Internet to record ever s from different information systems, and correlate and share these among the part. vs in order to detect and monitor frauds. A limited set of cyber threats belonging to inter-domain port scans are considered. An automated framework to respond to  $c_{b}$  the thr ats with appropriate response plans is proposed in [28]. It works by integrating and valuating operational, financial and threat impact models. The data emplyed is etwork inventory, security policies, mitigation actions, reachability matrix and vulner; bility inventory. The data is local to the network and based on preestablish d security guidelines. The proposed framework does not employ CTI data for dynan ic risk nanagement. In [29], a fog-based storage approach is proposed for the integri v, avai<sup>1</sup> bility and confidentiality of users data stored in clouds from a limited set of cyber timeats such as malicious modification and data loss.

Pub iciy available CTIR repositories such as VERIS incident database, Hackmageddon, . id the Web Hacking Incidents Database have been used by Liu et al [30]. The purpose is to forecast the cyber security incidents based on the network s<sup>+</sup>... This state is depicted with externally detected properties i.e mismanagement sy apto is, DNS or BGP misconfigurations and a time series of malicious activity including so uning, phishing and spam. The VERIS incident database and Alexa Web Information Service AWIS is used by Sarabi et al [31] to devise a categorization metric. It suggests the investment a business should do on achieving targeted and effective resource alle atom for security incidents dealing with critical data of that business. The incidents a categorized based on the type, the originating source, motive, and assets user' related to that incident. Similarly, Vasily [32] used the VERIS incident database to i vestiga  $\exists$  data breaches in financial sectors. To manage operational risk a belief netwo, ' is f rmed. The results obtained from the model assist management in underst .ndip~ the problem areas relevant to their businesses and make decisions for future tott rme t. In [33], the authors propose the use cases of cyber threat intelligence shaling usin, STIX to gain information about the disruptive consequences of cyber threats and data breaches on society. A threat analytics framework to contextualize the massive Cyber Threat Intelligence (CTI) is proposed in [34]. The CTI standards, n. twon. configurations, and Common Vulnerabilities Exposure (CVE) are represented usine ontological formal specification The ontology developed is employed to seman. Ily reason about the cyber threat relevance to a network their likelihood and the vulner, ble and affected assets. The focus is on the network relevance towards a cybe. the call instead of detecting or correlating existing threats with new cyber threats. In [5], the authors performed a comparative analysis of the STIX, IODEF, VERIS and A ALT cyber security incident reporting formats. The study reveals that STIX is the most comprehensive and practicable cyber security incident reporting format. In 30, the authors describe threat intelligence use cases with STIX. The purpose is to improve the automated management of threats. The information from different threat *t*elligence providers and network elements is correlated. A comprehensive survey of the existing CTI standards in the context of cyber threat information sharing is  $g_{1}$  in [7].

Security concepts extract d from unstructured online web sources proved to be really helpful in enhancing tradit onal knowledge bases of intrusion detection systems. More et al. [38] proposes an intru- on  $\alpha$  '-ec' on system which extracts security concepts from text. These concepts are cor, ared with monitoring sensors' logs with the help of a reasoner for generating security alere. The security concepts are extracted from heterogeneous sources. The extraction threat is used to populate a security ontology borrowed from [39]. In a related rese rch, Mulwad et al [40] present a framework for vulnerability extraction and cyber *tack* related information from web text and compare it with wikitology conc. pts [41]. A model proposed by Joshi et al [42] takes text, vulnerability descriptions, b. or osts and security bulletins as input and automatically extracts entities and populater concerns in DBpedia [43]. In DBpedia, these concepts are matched and assigned corresponding class values. Similarly Bridges et al. [44] present a maximum entropy mc lel for : utomatic labeling of security text. An approach for early detection of real-worl express and vulnerabilities from Twitter is proposed by Sabottke et al [45]. It observes tweet of security vendors, hackers, and administrators. More vulnerabilities can be detarted using Twitter security forums than actually present in the proof of concept public repositories. To extract useful content from news stories Ryan et al. [46] use a text mining approach. To measure relatedness between concepts collected from news stories Laten Semantic Analysis (LSA) is used. To semantically contextualize IoT data, Mário

et al. [47] proposes an approach for extracting semantic features from public 'y accessible web services to facilitate M2M based communication. A cyber security ontology called CRATELO is introduced by Oltramari et al. [48]. It consists of three is that use a simplified version of Dolce named Dolce spray as first level ontolog. The proposed cyber security related ontology SECO as a middle level and domain recel ontology OSCO for cyber operations. The research considers spatial and tempor 1 propercies of attacks and the attacker's host and network artifacts are not considered. In the related research the same authors [49] introduced the concept of trust management in CRATELO, e.g., any value delay which is out of an acceptable network delar range will be unreliable. Details of acceptable delay for different networks are not given. To it vestigate APTs via artifacts collected from cloud apps, Christian et al. [50] developed a taxonomy of security mechanism circumvention techniques by analyzing case. The concepts of the same security appears.

A Support Vector Machine (SVM) based signature free realware detection based on the selection of filter based malware feature is given in [51]. The features employed are Application Program Interface (API) call statistics. Similarly, a machine-learned model based malware detection system is provided in [52]. The malware features are obtained from both benign and malicious executables. The features considered are file hashes, malicious IP addresses and malicious externational statistics. The machine learning techniques used are SVM and decision trees. In [53], the author, propose a machine learning based malware detection engine. The malware features are on ained from APK files. The behavioral and permission related. The static features are on ained from APK files. The behavioral features are obtained by executing the Ap<sub>1</sub> sing an emulator. The permissions are selected using information gain. The J48 (open sourch Java-based implementation of the decision tree) classifiers outperforms other techniques. In [54], a machine learning based approach is used to detect and predict APT. The features of APT considered are low-level IOCs, i.e., executable files, malicious file instances and source and sources and source of the source of the decision to the detect of the static features are of the features are solved are low-level IOCs, i.e., executable files, malicious file instances and the source of the techniques are solved are low-level IOCs,

The aforementioned works is produced the groundwork for cyber security related information retrieval and set the stage for in estigating cyber data breach incidents. However, the models and ontologies er ployed and very basic and do not cover all aspects of threats both low-level artifacts and TT 's. They only capture vulnerabilities exploited and their needs. The proposed threat in stigation is based on semantic analysis of adversary's attack patterns by emploring machine learning techniques for predicting the threat family class. In order to identify an ats in the context of recent data breaches and exploits found in CTIR.

### 5. Data Brezch Juvestigation Framework

A high-level are it seture of the proposed framework is shown in figure 3. It can be broadly divided in to three segments: Semantic Indexer and Retrieval System (SIRS), TTD Semantic Ne work and Cyber Threat Prediction. The input to the system is CTIR and  $AT^{m}cCr$ . Locuments. A CTIR corresponds to a single cyber threat or incident while in AT'. CCK document may correspond to many detection mechanisms related to a 1 TP. The SIRS segment semantically indexes CTIR and ATT&CK documents, meintains a 1 TP dictionary extracted from threat documents and retrieves a ranked list of CTIP, and ATT&CK documents for each TTP present in the dictionary. The list of ranke, documents is combined to connect cyber threats, TTPs and detection mechanisms based on a higher rank specified by the ranking function. A belief net  $\ldots$  k is further trained between the TTPs and the cyber threat incident to predict cyl er tl reats based on detected threat artifacts. The last segment, i.e., *Cyber Threat Prea.* ion takes as input the *TTD* and the detected threat artifacts to produce a reliable  $\ldots$  reat prediction. The functional details of each segment are given in the following subsections. The details of notations used are given in Table 1.

#### 5.1. Semantic Indexer and Retrieval System

Table 1. Data Dreach Investigation

Notation	Definition
$A_{li}$	A ATT&CK documents set, associated with the TTP $ttp_i$
$TTP_{ti}$	A TTPs set, associated with a threat
$T_{si}$	A threats set, associated with be TTF $ttp_i$
$D_{di}$	A set of detection mechanis. $\circ$ . associated with $d_i$
TIndex	Latent Semantic Indexed threat Cocuments
AIndex	Latent Semantic Indexe. ALT&CK documents
TRank	A list of ranked threat documents associated with $ttp_i$
ARank	A list of ranked ATT $\sim$ 1 comments associated with $ttp_i$
TTPD	A detected ttps set
$TTPD_{ti}$	A set of detected $\sum$ due to threat $t_i$
$TTPU_{ti}$	A set of undetected $T_{1}$ due to threat $t_{i}$
$p_i$	A set of predicted threats based on $TTPD$
$S(t_i)$	Threat support function to measure the maximal support of
	TTPD towards $t_i$
P'	A new set $f$ threat predictions, $p_i$ based on $TTPD$
TTPM	A set of correlat d but missing TTPs associated with threat in
	a prediction
TTPE	A subset of $\square^{T} PM$ that are present
TTPN	A $\varepsilon$ abse' of $TTPM$ that are not present

The structure of a scheral CTIR document is shown in figure 4. In the case of this example, it is an excerpt of a STIX encoded CTIR document reporting the Backoff malware belonging to the mancial systems-breach threat family [55]. Each CTIR document is assigned a unique identification number using the *id* attribute in the *STIX\_Package* element. The low-level <sup>1</sup>OCs are represented by the *Indicator* STIX element. The high-level IOCs are represented by the *TTPs* STIX element. Both kinds of IOCs have unique IDs to disting the *in* the *TTP* element. The details of the TTPs are given in the *Description* sub-element of the *TTP* element. The STIX standard does not directly support the specification of detection mechanisms, which are typically not found in cyber threat reports. This is not one can be security administrators to identify TTPs in their network, those ound not the CTIR must be mapped to their associated detection mechanisms. The detection mechanisms are provided in the ATT&CK repository by MITRE [21]. In this repository, ten tactic categories are defined i.e. *Persistence, Privilege Escalation, Defence Exactor*.



*Exfiltration* and *Command and Control*. Each category specifies the advantsary's techniques with threat examples and actions for detection and mitigation. An example of an ATT&CK document describing *Application Deployment Software* technique's is shown in figure 5. Each technique has its own unique identification number. The 'echnique's title is a generalized taxonomic label which is difficult to map to complex insidely described TTPs in the CTIR.

The SIRS algorithm is provided in listing 1 and it's detailed 'xplanation is provided here. In addition, its working with respect to a sp cinc case study of financial malware is provided in section 6.2.1. The SIRS segment constructes a main function SEM\_INDEX\_RETRIEVE whose job is to retrieve ran. d C'. IR and ATT&CK documents for each TTP. The function takes as input a set of m CTIR documents  $T=t_1, t_2, ..., t_m$  and a set of l ATT&CK documents  $A=a_1, a_2, ..., c$ . T represents individual attack instances, malware families or software tool. A set  $\circ^{c} k$  detection mechanisms  $D=d_1, d_2, ..., d_k$  can be used to verify existence of a set of n TPs  $TTP=ttp_1, ttp_2, ..., ttp_n$ . The first step, shown in line number 2 to 4 extracts TTPs f om the CTIR and archives them in a dictionary. As mentioned in the introdu *tion* ....tion, we have used STIX encoded CTIR documents. STIX uses the commonly use, 'XML format to import, export and share data between different platforms. I. vever, to analyze XML encoded data, these files can be converted into different flexible for ats. A common practice is to import the XML encoded data into spreadsheet ', 'I chiract TTPs from the STIX encoded CTIR, a dictionary is created on line number - whereby the CTIR files are imported as separate rows into a spreadsheet and the continue constitute the CTIR ID, TTP ID and TTP Description. An excerpt of the data s'ructure of the TTP dictionary is shown in figure 6.



Figu 9 4: An xcerpt of STIX encoded CTIR document reporting Backoff financial malware [55]

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Figure 5: An ATT&CK MITRE document specifying adverse v's TTP, threat example, mitigation and detection mechanisms

CTIR ID	ТТР Ь	TTP Description
Feed- 4170e5e758e50d29934ad e2e19cd0c02	TTP-b500e، م-4955 کی ،3- d8j2-8c3d4b304	UsesTOR hidden services
Feed- 4170e5e758e50d29934ad e2e19cd0c02	T ,-f9d6e。 >-a7d6-0094- 7%c-535f8c(:)1d73	Encrypts every file it can find that isn't an EXE or DLL
Feed- 4170e5e758e50d29¢,4ad e2e19cd0c02	T <sup>*</sup> -006e0a14-677b-dfed- ,447-4 59298c6272	Downloads the encryption key from its CnC server

r. vre 6: An Excerpt of TTP Dictionary

The second ster shorn on line number 6, is to semantically correlate the TTPs in the CTIR among themsel. s (line number 7) and with the TTPs in the ATT&CK documents (line number 8) In figure 7, the semantic relevance among TTPs in ATT&CK documents and CTIR is boom 1. Con the left-hand side, there are five ATT&CK documents. The document's title is be table of the TTP. On the right hand side, there are three different threats of manciel malware family, i.e., Backoff, Treasure Hunt and FastPOS. It can be seen that there are some common TTPs, e.g., key logging functionality and Deliver keylog  $d \uparrow^* r$  buching the textual descriptions, it is hard for a machine to correlate them rising a limple keyword matching technique. Similarly, It can be seen that the titles of rele bant AT f&CK TTP documents represent the general labels or taxonomic classes of the Times, which is again hard to correlate with the textual TTP description in the Cin 'R. For this purpose, we use LSI to semantically map TTPs with each other. In figure 7, the tTPs are numbered on both sides. This numbering represents their semantic relevance to the semantic relevance of the transmanner of the titles.

vance which is the desired outcome of using LSI. In the past, researchers b  $\therefore$  successfully used LSI for enhancing the semantic understanding of text documents in web search [56], recommendation systems [57], personality identification [58] and automa  $\therefore$  assignment of multi-level security labels [59]. The traditional keyword-based lexical matching techniques retrieve information by matching query terms with the term  $\Rightarrow$  in the text corpora. However, they can be erroneous matches due to inaccurate concrept matching caused by synonyms and polysemous words. Thus instead of using the simple retword search, we index CTIR and ATT&CK documents using LSI. The TTPs are searched for semantically relevant concepts or topics in statistically derived conceptual in lices (line number 11-24).

LSI assumes a latent pattern in document terms which is biolen due to variable ways to represent a concept. The dimensionality of the sparse (with r ore zero entries) termdocument matrix (representing the frequency of termediate of the documents) is reduced using Singular Value Decomposition (SVD) to figure out the hidden word patterns in the documents. The relevance between the TTP query and the indexed documents is calculated by taking the cosine of the two vectors. The documents are ranked by the distance to the TTP queries. A smaller angle leads to a longe cosine value that means the document has high relevance to the given TTP query. A larger angle leads to a smaller cosine value that means the document is low in relevance to the given TTP query.

To connect TTPs in the dictionary among 'her ascives, the TTP descriptions in column 3 of the dictionary are provided one by one a search terms in the CTIR corpora. If there is a match then a new row is created in the dictionary with the same TTP ID as the ID of the TTP that was matched. Similarly, to connect TTPs in the dictionary with the detection mechanisms, the affect accorptions are provided as search terms in the ATT&CK corpora. If a match is found the corresponding detection mechanism is fetched and added as a new columer in the dictionary.

The Latent semantic indexe<sup>-</sup> indexe<sup>-</sup> CTIR and ATT&CK documents. The Semantic retrieval system retrieves a ranke<sup>-1</sup> list *c* i CTIR and ATT&CK documents for each TTP present in the TTP diction<sup>-1</sup>ry. The lanked results against each TTP are merged and further provided to the sec<sup>-1</sup> d egm<sup>-1</sup>nt, i.e., TTD network where TTPs are semantically linked to their corresponding enterlated in mechanisms.



Figure 7: Semantic correlation among TTPs . C<sup>TI</sup>R and with TTPs of ATT&CK documents

### Algorithm 1 Semantic Indexer and Rev. oval System

```
Input: CTIR T, ATT&CK documents A
Output: An n-element array of Ra. ked CTIR and ATT&CK documents (TRank,
    ARank) for each TTP
 1: function SEM_INDEX_R'_TRIEV_'(I', A)
        TTP_Dictionary = F_{xtre} t_Relevant_Fields(T) > Populate TTP_Dictionary from
 2:
    CTIR
        LSI_Indexer(T, )
 3:
        Semantic_Retri val "TP_Dictionary, TIndex, AIndex)
 4:
 5: end function
 6: function LSL' NDE ER(T, A)
        TIndex=LS1
 7:
        AIndex= \angleSI(A)
 8:
 9: end funct on
10: function SEN, 'NT.C_RETRIEVAL(TTP_Dictionary, TIndex, AIndex)
        for = 1 to m do
                                             \triangleright m = count(ttp<sub>i</sub>) where ttp<sub>1</sub>, ttp<sub>2</sub>, ...., ttp<sub>m</sub> \in
11:
    TTP_I ictionar
            Sumplies tp_i in (TIndex, AIndex)
12:
            TI = \operatorname{Rank}(TIndex) for ttp_i
                                                      ▷ Where Rank - LSI document ranking
13:
            Self et tl_i \in TL_i > Threshold
14:
            f_{-}, j = 1 to n do
15:
                                               \triangleright n = count(tl<sub>i</sub>) where tl_1, tl_2, \dots, tl_n \in TL_i
16.
               TRank = add_Record(tl_i, ttp_i)
            end for
17:
            AL_i = \operatorname{Rank}(AIndex) for ttp_i
18:
                                               14
            Select al_i \in AL_i > Threshold
19:
            for k = 1 to r do
                                              \triangleright r = count(al<sub>k</sub>) where al_1, al_2, \dots, al_r \in AL_i
20:
21:
               ARank = add_Record(al_k, ttp_i)
            end for
22:
23:
        end for
        return TRank, ARank
24:
25: end function
```

### 5.2. TTD Semantic Network



The SIRS segment retrieves semantically relevant cyber threat report and adversary's techniques documents. The next step is to link threats to their respect  $\gamma$  TTPs and detection mechanisms. A dependency table  $P_{T\times S} = p(ttp_i/t_i)$  is built betoen the threat incidents and the TTPs. Thus TTD represents semantic relations of bree independent concept sets, i.e., threat set, TTP set and detection mechanisms set. fTr s and threats are connected through links to show that the two are dependent. TPs and detection mechanisms are also similarly connected to show dependencies between the TTPs and the detection mechanisms as described in the ATT&CK taxe nomy.

Algorithm 2 Threat-TTP-Detection Semantic Network
Input: TRank, ARank
Output: TTD Semantic Network
1: function TTD_Semnet(TRank, ARank)
2: TTD_NETWORK $(ttp_i, TRank, ARank)$
3: BELIEF_NETWORK( <i>TTD Network</i> )
4: end function
5: function TTD_NETWORK $(ttp_i, TRank, Ahh^k)$
6: for $(ttp_i \in TTP)$ do
7: add $(t_i \in TRank)$
8: add $(d_i \in D \in ARank)$
9: end for
10: return TTD Network
11: end function
12: function Belief_Network(TTD Network)
13: for $t_i \in T$ do
14: $\omega(ttp_i/t_i) = \frac{p(tt_i/t_i)}{\sum_{t_i \in T_{ttp_i}} \gamma(ttp_i/t_i)}$
15: $\mu(t_i/ttp_i) = \frac{(ttp_i/t_i)_{t_i}}{\sum_{t_i \in T_{ttr_i}} \omega(ttp_i/t_i)p(t_i)}$
16: end for
17: return TTD Se vantic Network
18: end function

The stepwise details to build a TTD semantic network are given in algorithm 2. The working of the TTD is explained in section 6.2.2 with the financial malware case study. After initial mapping between threats, TTPs and detection mechanisms, a belief network is trained betworn threads and TTPs in order to predict threats based on the presence of certain artifacts. It is based on a threat support function  $S(t_i)$  to measure the maximal support of the det sted TTPs towards a threat occurrence. The algorithm assumes that all predicts TTPs are independent of each other. The historical artifacts signifying the existance of a constraint of a constraint of the data threat are gathered to calculate conditional probability between TTPs and the sats, i.e.,  $p(ttp_i/t_i) \in (0, 1)$ . These probabilities are based on historical data that form is a frequency table in the Threat-TTP mapping. The frequency table may need to be normalized to eliminate null values. This is done by adding 1 to all entities of the apping table and thus eliminating entries with "0" or null value. The normalized to calculate the normalized likelihood or normalized conditional

probability  $\omega(ttp_i/t_i)$  to show the support of each threat  $t_i$  to their determinant  $t_i$  shown in equation 1.

$$\omega(ttp_i/t_i) = \frac{p(ttp_i/t_i)}{\sum_{t_i \in T_{ttp_i}} p(ttp_i/t_i)} \tag{1}$$

Using the above normalized likelihood table, the normalize ' pc sterior probability  $\mu(t_i/ttp_i)$  can be computed by using Naive Bayes in equation 2 [22]:

$$\mu(t_i/ttp_i) = \frac{\omega(ttp_i/t_i)p(t_i)}{\sum_{t_i \in T_{ttp_i}} \omega(ttp_i/t_i)^{p(t_i)}}$$
(2)

The predicted class is the one that has the highest proterior probability. All the detected TTPs are considered for such support and for each attack for a different kind so in order to extract a threat prediction all such support values are combined to find a best candidate threat prediction set with maximum  $\sup_{\mathbf{k}} \operatorname{pred} x$  value  $S(t_i)$  shown in equation 3.

$$S(t_i) = \frac{\sum_{ttp_i \in TTPD_i} \mu(v_i, {}^{\prime}ttp_i)}{\sum_{ttp_i \in TTP_{t_i}} u(t_i/ttp_i)}$$
(3)

The output of the second segment is a TTP network which is provided as input to the third segment, i.e., cyber threat predictio.

#### 5.3. Cyber Threat Prediction

The cyber threat prediction segment can be further divided into three functional modules: 1) Threat Investigation (TD), 2, Reliability Assessment (RA) and 3) Detection Mechanism Selection (DMS). The stepwise details of the cyber threat prediction segment are given in algorithm 3. The working of each functional module is explained in section 6.2.3 with POS malware case story.

The responsibility of the ' hreat  $\dots$ '' stigation module is to produce a predicted threat set P' given a set of detecte  $_{i}$  TTPs i e TTPD. P' may contain a set of attack predictions  $(p_1, p_2, \dots, p_n)$ .  $p_i$  is a threat  $\neg$  t that characterizes the TTPs detected. Next, the reliability module determined if any of the predictions  $p_i(p_i \in P')$  is reliable. The first reason for reliability assessment is that a prediction with low reliability can lead to an incorrect prediction and thus waste the time and resources of an organization. The second reason for assessing reliability  $\cdot$  to determine the presence of poisoned or spurious symptoms, i.e., detected TTPs. The Reliability Assessment (RA) equation used to measure the reliability of prediction  $_{F_{i}}$  is given in equation 4. The  $TTPD_{ti}$  represents a set of all detected TTPs due to threat  $t_i$ . The  $TTP_{ti}$  represents a set of all TTPs, associated with threat  $t_i$ .

$$RA(P') = \frac{\sum_{t_i \in p_i} (TTPD_{ti}/TTP_{ti})}{p_i} \tag{4}$$

The "real investigation is deemed complete if a high reliability is determined. If the relevant TPs related to the threat prediction  $p_i$  were not detected then, *TTPM*, the set of und dected TTPs that most reliably contribute to explaining  $p_i$  are considered by the Detection Mechanism Selection module. This outputs a set of existing TTPs, i.e.,  $TT \sim and$  non-existing TTPs TTPN based on the least cost detection mechanisms and t.  $\circ$  prediction  $p_i$  reliability value is recalculated. If the reliability value increases

```
Algorithm 3 Cyber Threat Prediction
Input: TTD Network, TTPD
Output: Reliable Prediction
 1: function THREAT_PREDICTION(TTD Network, T, P)
       THREAT_INVESTIGATION(ttp<sub>i</sub>, TTD Network, 1 PD)
 2:
 3:
       RELIABILITY_ASSESSMENT(P')
 4: end function
 5: function THREAT_INVESTIGATION(TTD N-1, PTD)
                \sum_{ttp_i \in TTPD_i} \mu(t_i/ttp_i)
       S(t_i)
 6:
                \sum_{ttp_i \in TTP_t_i} \overline{\mu(t_i/ttp_i)}
 7:
       P' = \operatorname{Max}(St_i)
       return P'
 8:
 9: end function
10: function RELIABILITY_ASSESSM. (\mathcal{D}')
                         (TTPD/TTP)
       RA(P')
11:
                           p_i
       return P'
12:
13: end function
14: if P' > Threshold then
       Reliable Prediction
15:
16: else
       if TTPMexists ' nen
17:
           Threat_Inversion (TTD, TTPE)
18:
       else
19:
           select lov cos d_i
20:
           Reliabil. \neg A sessment(P' = TTPE, TTPN)
21:
       end if
22:
23: end if
```

sufficiently then the prediction procedure terminates, otherwise the three investigation module receives TTPD, TTPU for a new threat prediction.

### 6. Financial Malware Case Study

The working of the data breach investigation framework is ellow the vector with the case study of financial memory scrapping malware [60]. The malwal steals identity and payment records by reading device memory at the retail chlockout. The TTPs of a few malware instances extracted from multiple sources [24] [15], [3] are mapped to their appropriate cyber kill-chain stages to understand the data dueft process. The TTD map is built by semantically mapping TTPs extracted from multiple sources to financial malware classes and their corresponding dete tion mechanisms as defined in the ATT&CK repository. The TTD network and the detected due artifacts are given as input to the cyber threat prediction segment to invelving threat occurrence and the presence of more TTPs by applying appropriate detection mechanisms. A detailed working is explained in the following subsections.

Reconaissance	econaissance Weaponize & Deliver		Insta	Insta Command and Control	
Phishing	Hardcoded CnC addresses	Handle Ser- vice Cuuu I Request	Use Enum Pro-	Use HTTP POST	Collect System Information
Drive by Download	Hardcoded Process Names	Use I hole 32Snaps. t method	Custom Search Function	Use HTTP GET	Injects Code
Luring Attack	Hardcoded File Names	U – Blacklist	Use Regexes	Use FTP Server	Exfiltrates data as Hex digits
	BOT Functior	1. ~ks API	Encodes data as Base64	Use SMTP	Manually data exfiltrated
	Socially In- gineered Fr. names		Encodes data as RC4	Use TOR	
	Preten s he JAVA Cod in VB- Scipt Culla (abfusca- Jion		Store Results in logs Performs Luhn Validation Installs Watch- dog RDP Brute Force Source Code Updated Exploit Auto Start Runkey Copy files in %APPDATA% Registers to by- pass firewal Has a Kill Switch	Use Interne- tOpen URL Send PHP() subroutine Grab Browser Form	

Table 2: Financial Malware TTP mapping to Cyber Kill Chain Phases

#### 6.1. Mapping Financial Malware to Cyber Kill Chain Phases

Different TTPs have been used over the ages to steal payment  $\epsilon$  rds' data. One technique is to target victims by downloading keylogging tools on their stems. But unfortunately, keyloggers are not capable of retrieving data stored on ca. 's' tracks. The attackers concentrated on victimizing corporations rather than in available that employ POS terminals to handle transactions conducted through cards The vulnerability targeted involves stealing cards' information temporarily held in the . omory of point of sale terminals as clear text. Therefore, a specialized malware is developed for harvesting the data stored in the memory. This specialized malware s typically called a 'RAM for stealing payment card details in various ways. To a curation diagnose the family of RAM scraper, the security analyst has to perform certa. investigative actions to discover the behavior and TTPs of the malware. Thus the "TPs are mapped to the phases of the cyber kill chain model to elaborate the theft incide. TTPs of eleven financial malware samples, i.e., Rdasrv, Alina, VSkimmer, De. 'er, B'ackPoS, Decebal, JackPOS, Soraya, Chewbacca, BrutPOS, and Backoff are concrete and compared. The TTPs for each cyber kill chain phase are discussed below with a mmary provided in Table 2.

In the *reconnaissance* phase, phishing techniques are often adapted to break into the victim's account in a network. Other techniques used are to discover the structure of the network and their exploitable vulnerabilities. The purpose of the adversary in this phase is to identify the victims and vulnerable resources of the target network.

The detection mechanisms related to this phase comprise all those proactive measures that should be taken to avoid attackers giving a chance of identifying network structures and selecting vulnerable targets. This is cludes user training, enforcing the principle of least privilege for all the users at every access and constantly assessing the network for risks associated with each vulner point v [61].

Once the intruder gets to know about his target, the malicious content is prepared based on the type of vulnerabilities present. This phase is known as *Weaponize & Deliver*. Once the malware is prepared it is sent to the target network bypassing the security controls. Malware is forn, that draget processes to be searched in the RAM of the POS device are hardcoded in the binary. The file names used to store malware on the system are also sometimes hardcoded in the binary. The file names are carefully selected using social engineering to the quest to avoid detection. In the same way, the Command and Control (CnC) address is to transfer the stolen content are also hardcoded. Again the investigation doe in this point. It starts when the malware is finally delivered to the target. So these two phases also depend on the proactive measures taken by the target as discuss of in the previous step.

The next phase is the *exploit* phase in which the malware actually begins to exploit the vulner bilities. The TTPs used in this phase register the POS malware as a service. The execution of the malware is controlled by the service control handler. To detect this TTT, the security analyst performs certain detection mechanisms. In the first step, chang s to the service registry are monitored. Secondly, the frequency of modification to existing services is checked. Thirdly, the security analyst looks for any abnormal change to invice binary path location which is not conventionally meant for that service and is also irr levant to software updates. After successful registration, the POS malware starts

scanning the RAM. The malware uses the *CreateToolhelp32Snapshot* meth  $\gtrsim$  for scanning the active RAM processes. A list of blacklisted processes can help in *i* lent; ying tracks 1 and 2 card data. Some malware uses regular expressions to do this,  $\downarrow$  it is a slow process. An efficient way for scanning processes is "process enumeration, in which all the processes in the POS RAM are scanned to search track 1 and 2 dat. The adversary uses the Windows Management Instrumentation (WMI) [62] tool for process enumeration. To detect process enumeration, the security analyst should monitor process executions. To obfuscate the malware's extracted content, the adversary use *Base64* encoding [21].

The next phase is the *installation* phase in which the ma ware trips to be persistent. There are multiple steps followed to achieve persistence, e.g., where code may be actively developed and regularly updated by installing latest code versions from the Internet during execution. It can be detected by carefully examining triggers from antivirus and other security tools. The malware may add a file path the autopoint of art run key that boots it each time the system is turned on. This can be checked by detecting irrelevant changes to registry run keys and checking the start folder here changes and additions. Similarly, the malware may add itself to a firewall as an autopoint of an equipplication via registry keys or install a keylogger in the current working directory. This can be detected by monitoring irrelevant changes made to the accessibility utility binaries or binary paths. Keylogger installation can be tracked through changes made to "DLLs [21].

The next phase is the command and cont.  $vt p_{1acc}$  in which the malware establishes a covert connection with the CnC server to give the attacker a direct control over the target network. This connection may be established a using HTTP POST/ GET requests, connecting to FTP servers, using emails throug. SMTP protocols and bots. To detect CnC communication the detection mechanic is used by the security analyst include analyzing network data for uncommon data flows via unseen processes, processes communication layer data. Similarly, the security analyst can monitor the processes for their file access patterns and network behavior. To the equest.

Action and objectives i the fina' cyber-kill chain phase in which the attacker steals payment card data, ten porally cores it on the disk and then exfiltrates that data. To collect system information, host enumeration techniques are used. To detect host enumeration, the security and 'vst investigates the execution of processes and detect the time when programs and doing system enumeration. The presence of keyloggers can be detected by observing ruodification in the DLLs. To detect code injection, the process registries should be an 'vzed for changes. Similarly, data encryption can be detected by finding high data encryption [21].

Thus we have seen not for each phase of the intrusion kill chain model there exist adversary patterns in the form of TTPs that can be mapped to their particular attack and detection mechanisms.

### 6.2. The st Diagnosis for Financial Ram Scrapper Malware

The works g of the data breach investigation framework is detailed here with the help of case study of financial RAM scraping malware. The details of each segment are given in the following subsections.

### 6.2.1. SIRS for financial Malware

The process of threat diagnosis for data breach incidents caused by f nan al malware starts with building a TTD network by extracting TTPs from financial n. ware related CTIR and semantically indexing them using LSI. Once when all the T'1. are extracted, semantic relevance among every single TTP and all the financial 'aa vare astacks mentioned in CTIR are found by searching the TTP in the LSI ind xed mancial malware CTIR. In figure 8, the search result of financial malware TTP re. +ed to establishing command and control connection using email or SMTP prote of is shown. The keyword "use email or SMTP for exfiltration" is searched in semant cally in lexed CTIR documents and we get relevance probability for each. Based on the . sult it can be seen that the TTP is present in BlackPOS, Dexter, and Vskimme . Th: means that the TTP is relevance can be mapped in the form of a Threat-TTP ... eque table as shown in Table 3 where the table entries specify the frequency of threat in ident occurrence. Similarly, we indexed all the cyber attack techniques mentione. in AT<sup>7</sup> &CK framework using LSI. Each TTP is searched against the indexed docume. 's to nok for the relevant detection mechanisms. In figure 9, the search result of financial nalware TTP related to establishing command and control connection using mail or SMTP is shown. The keyword "use email or SMTP for exfiltration" was searched in semantically indexed ATT&CK documents and we get relevance probability for each.



Figure 8: Search results for "use email or SMTP for exfiltration" TTP in LSI indexed CTIR

Table 3: Financial Malware Threat-TTP Frequency Table							
TTP ID	Description	RDASRV	Alina	Vskimmer	Dexu	BlackPOS	
TTP1	Credential Compromise	2	4	3	F	5	
TTP2	Known POS System	2	1	1	1	1	
TTP3	Hardcoded Filenames	1	5	3	2	1	
TTP4	Changes registry as le- gitimate Firewall App	1	1	5		1	
TTP5	Hooks WH_KEY BOARD_LL	1	1	1	5	1	
TTP6	CreateToolHelp 31 Snapshot Method	1	2	2	F	1	
TTP7	Use Email or SMTP	1	1	3	6	2	



Based on the results it on be seen that most relevant detection mechanisms related to the mentioned TTP are present in ATT&CK document termed as "Email collection". The detection mechanic in from this document are enumerated as follows:

- 1. Check for process utilizing network communication that should not require this function
- 2. Check for proce es accessing network without user-driven events

### 6.2.2. TTD N +wc k fc Financial Malware

After retrieving the lantically relevant financial malware related CTIR and ATT&CK documents a TT.) network is built shown in Table 4. The associated detection mechanism details are liven in Table 5. Here five malware instances are considered, i.e.,  $RDASP^{V}$  A<sub>thener</sub>, Vskimmer, Dexter, and BlackPOS. A belief network between threats and T iPs is formed by calculating posterior probabilities based on the incident frequency. The losterior and prior probabilities are shown in Table 6. For a predictor attribute, the financial malware class, TTP, that boasts the highest posterior probability is considered the prediction output. For instance, "Alina" has the highest occurrence probabilities for the malware threat class, and predictor

attribute, i.e., TTP before seeing any real-world data are also given. It c<sup>-</sup>... be seen that "Dexter" has the highest prior probability which is also evident from 'listor cal data of financial malware incidents. Similarly, it can be seen that among all  $1^{-1}$  Ps, TTP1 is the most prevalent TTP as its prior probability is high. The TTD network predicts the threat incident based on the detected threat symptoms provided in the form of host and network artifacts.

Table 4: Financial Malware Threat-TTP-Detection Ser antic Notwork							
Threat	TTP ID	Deter ion					
		Meclanism					
RDASRV, Alina, Vskimmer, Dexter	TTP1	41					
RDASRV	TTP2	.2, d3					
Alina, Vskimmer, Dexter	TTP3	d4, d5, d6					
Vskimmer	TTP4	d7, d8					
Dexter	TTP5	d9					
Alina, Vskimmer, Dexter	TTı	d10					
Vskimmer, Dexter, BlackPOS	'1."P7	d11, d12					
RDASRV, Alina, Vskimmer, Dexter RDASRV Alina, Vskimmer, Dexter Vskimmer Dexter Alina, Vskimmer, Dexter Vskimmer, Dexter, BlackPOS	TTP1 TTP2 TTP3 TTP4 TTP5 TT1 '1, "P7	11 12, d3 d4, d5, d6 d7, d8 d9 d10 d11, d12					

.

	Table 5: Detection Mechanish, 3 104,ncial Malware TTP	
Detection Mechanism	Description	Weight
d1	Look for suspicious account ben vior	5
d2	Audit file system aclogs _`r failed attempts	2
d3	Monitor process capture of the arguments	4
d4	Analyze file read, write and modify method for malicious ac-	2
d5	Monitor anti- rrus a. ' IDS alerts	1
d6	Acquire IP $\epsilon$ d port	1
d7	Monitor change. in ut lity arguments and binaries	4
d8	Monitor command in e invocations modifying registry	3
d9	Monitor regis ry for key stroke interception driver installation	3
d10	Monitor <sub>F</sub> _ ram or process enumeration	4
d11	Chec' for out ' context network communication	5
d12	Che retwork access without user driven event	5

	<u>Table 6: 1.</u>	n <u>cial Malwa</u>	re Threat-TTP	Posterior a	nd Prior Probab	ilities
TTP ID	R' ASPV	Alina	Vskimmer	Dexter	BlackPOS	Prior (TTP)
TTP1	0.1	0.2	0.15	0.3	0.25	0.26
TTP2	0.33	0.17	0.17	0.17	0.17	0.08
TTP	0.08	0.42	0.25	0.12	0.08	0.15
TTP4	0.11	0.11	0.56	0.11	0.11	0.12
TT 5	U I	0.14	0.14	0.43	0.14	0.09
′ .′TP6	ി.09	0.18	0.18	0.45	0.09	0.14
TP7	.08	0.08	0.23	0.46	0.15	0.17
Prio.	0.12	0.19	0.23	0.31	0.15	

#### 6.2.3. Financial Malware Threat Prediction

After generating the TTD network, the process of cyber threat in estisation starts with threat prediction. This happens when the network administrator of erves certain artifacts. The artifacts, that can be seen, are low-level indicators su<sup>h</sup> as IPs, CnC servers, and ports or it can be unseen high-level indicators such as 1. Ps. '1 nese indicators are either obtained directly or through certain detection m $\epsilon$  har sms. Suppose the network administrator observes a suspicious account behavior, cha. es in registry with the installation of a driver intercepting keystrokes and a process accessing the network without any user-driven event. Based on the derived TTD semantic etwork, the detection mechanisms selected are d1, d9, and d12. These detection  $mec^1$  anisms are mapped to TTP1, TTP5, and TTP7 which become our detecte. TTP- i.e., TTPD. Based on the TTPD, the belief network predicts the probability f ccur ence for each malware instance as shown in figure 10. It can be seen that "L'ockPOC" has the highest probability of occurrence. The second highest probability is that of "Dexter". Now we have to select a threat prediction set P'. A problem that "rises concerns determining exactly which malware instances should be included in the prediction set P'. For this purpose, a threshold must be defined based on expert opinion - ;, suppose for this example we set threshold as above as 20%. Thus the prediction set comprises two threat instances, i.e., "BlackPOS and Dexter". Setting a high thresh. Id leads to a more reliable threat prediction as a low threshold compromises t. 9 remainlity of the threat prediction.



Fig. e 10 Probability Prediction for Financial Malware Instances

We assess the reliability of the prediction using equation 4 given in the design section. In this case, the total rumber of TTPs for the BlackPOS malware are "2" and that for Dexter are "5". The detected TTPs, i.e., TTPD for BlackPOS are "2" and that for Dexter are "3". There are two malware in the threat prediction set P'. After inserting the values in equation 4, we get equation 5

$$RA(P') = \frac{\sum_{t_i \in p_i} (2/2) + (3/5)}{2}$$
(5)

Thus the reliability of this prediction set becomes 80% which is quite high. This sho is that the prediction is reliable. The reliability of threat prediction increases more if we herease the threshold for threat selection in the prediction set, e.g., we include

all those malware in the threat prediction set who has a prediction greases than 25%, then there will be only one malware selected, i.e., BlackPOS. The reliability for this prediction set is a 100% greater. On the other hand, if the threshold for set cting threats in the prediction set is decreased the reliability also decreases, e.g., we below all those malware in the prediction set that have a prediction greater than any then there will be three malware selected, i.e., BlackPOS, Dexter, and RDASRV. The reliability of this prediction set is 70% less than the previous prediction set.

Suppose that the reliability assessment of the threat production is low. Then the undetected TTPs of the malware in the prediction set, i.e., 'TPM r ust be considered. To check the existence of TTPM, the security administrator  $_{1}$  or for instheir associated detection mechanisms. This may be hard to accomplish if the promber of TTPs in TTPM is large. Thus, there is a need to find the least cost detec for mechanisms that can verify the undetected TTPs, i.e., TTPM of the threat pred. for writh a high reliability. One solution is to let the security administrator decide which content mechanisms should be performed based on his knowledge, expertise, entironment and the resources. In the capacity of this research work, we have assigned work the security actions based on the security expert opinion. The complexity of detection mechanism are assigned weights based on their complexity as shown in Table 5. The weights is in the range [1-5]. The detection mechanisms with a weight closer to '0' are considered less costly than those with a weight near to '5'.

For example, one of the malware's  $1^{+}$  is 10 use RC4 to encode the contents for exfiltration. The detection mechanisms related to this TTP are:

- 1. Use samples to acquire the key and it algorithm and employ the same for decrypting network traffic for determining the communications signatures for malware.
- 2. Inspect SSL/ TLS to determine h encrypted traffic contains the presence of a C&C or in other words a Commond and Control channel.

It is clear that detection .nec'.anisms for getting the encryption algorithm and its key is a more difficult process a compared to inspecting SSL and TLS. Thus the weight of SSL/TLS inspection is 'ow as compared to finding the encryption key.

#### 7. Evaluation and south

The effectiveness and efficiency of all the segments of the data breach investigation framework were evaluated.

The effecti ene s of SIRS was evaluated using the financial malware TTP dataset compiled in the care oudy using the parameters f-measure, recall and precision. The effectiveness of the 1 fD network was evaluated by constructing a benchmark training dataset from the threat, TTP and detection mechanisms data available in the ATT&CK taxonomy [2.] This benchmark data serves as a ground truth regarding cyber threats associated with different threat families for training our model by building a TTD network and provides a baseline for future comparisons. The TTD network is evaluated using a hold-out threat of detection and the false discovery rate. The accuracy of the threat detection procent was monitored for two kinds of situations. The first situation was when certain threat artifacts were lost. The second situation was when certain spurious  $\therefore$  eat artifacts were added to poison the threat investigation process. The effective ess i the cyber threat prediction segment was evaluated through the factors that impact the reliability of the prediction. These factors include varying levels of probability  $\vdots$  resholds of the prediction set and the number of detected artifacts and types of  $T^{T}_{1,c}$  (either disjoint or overlapping) controlled by attack type selection. Whereby having the same attack family provides more overlapping TTPs and using data about more rec. t attacks provides disjoint TTPs.

The efficiency was evaluated based on the time taken to detect an attack for different dataset sizes and threat records. The TTP coverage  $n^{-1}$  the threat incidents of the ATT&CK dataset and the capability of our protot pe statem in data augmentation is discussed in the light of well-known threat incidents related to particular threat actors reported in multiple sources. Finally, a feature based comparative analysis of the threat investigation framework was performed with exacting intrusion and attack detection techniques to elaborate on the benefits of the proposed approach to the overall security measurement and management infrastructive.

#### 7.1. Benchmark Dataset

To construct a benchmark dataset, ground truth data related to cyber threat incidents, encompassing the details of TTPs and delection mechanisms, was required. For this purpose, we conducted a thorough murvey of literature available on standardized cyber threat taxonomies. The three most we'l-known are the Open Threat Taxonomy [64], the Enisa Threat Taxonomy [65] and the ATT&CK Taxonomy [21]. Among these, we found that the Open Threat and Enile did not serve our purpose because they simply list threats along with their descriptions but do not distinguish the TTPs and detection mechanisms. On the other hand and the threat threat incidents and connects them to the TTPs and detection mechanisms, and is, therefore, the closest to the requirements of our training model.

TTP Classes	Doc ments	hreat Actors	Documents Software Tools		Documents
Persistence	28	Chinese	17	Malware	65
Privilage Escala-	7.	Russian	3	Backdoor	37
tion					
Defense Evasion	23	Iranian	2	Bot	3
Credential Access	8	Indian	2	Web Shells	3
Discovery	16	Portuguese	1	Toolkit,	4
				Bootkit	
Lateral Movement	14	Unattributed	20	Rootkit	4
Execution	1.			Credential	4
				Dumper	
Collection	11			Credential	1
				Harvester	
Exfilt <sup>,</sup> .tion	9			Remote Access	9
				Tools	
Comm. rd and	17			System Utili-	16
Control				ties	

Tably 7: . T& K Taxonomy Dataset Statistics [21]



Figure 11: TTP distribution for Throat Incident



Figure 12: Three in tent distribution for ATT&CK TTPs

To establish the reliability of the AT f&CK data we investigated its sources. "ATT&CK" is created and maintained 'y MATRE Corporation [21]. It covers adversary's TTPs encountered in pre-attack a we, as ost-attack stages proposed by Lockheed Martin [66].

At the time of writing there exist 133 adversary's technique documents classified under 10 tactic classes. At tech, 'que may belong to more than one class. We have compiled some key statistics for the interest of the reader in Table 7. Each technique document has a short description alone, with its mitigation, detection mechanism, and threat incident example. Threat incident examples are either threat actors or software tools used by them. At the time of writing, there are 45 threat actors and 123 software tools in the ATT&CK tax nor y. Figure 11 shows the TTP distribution of threat incidents. The minimum number of TPS encountered in a threat incident is 1 and the maximum is 34. An average TTP count for each threat incident is 6. Similarly, figure 12 shows the other side of the picture, i.e., threat distribution for TTPs. The range of TTPs lies between 0 and 49. An average threat incident count for each TTP is 7.

The ATT. CK TTP taxonomy at any snapshot in time does not always encompass all TTPs of ary given threat mainly because the website is constantly being updated as more into mation about a threat is uncovered. One can observe cases of threats where convections are obviously missing between threat incidents and TTPs. For instance, consider threat incidents caused by a hacking group known as Axiom or Hidden Lynx [67] extracted from ATT&CK. This group offers hacker hirity services and is known to have launched highly targeted campaigns against diverse business domains including multiple sectors and can run multiple campaigns simultaneously. It is observed that only three TTPs are listed for this group in the ATT&CK repository as shown in Table 8 while there are several more TTPs that can be related to this group obtained from other sources such as IBM X-Force STIX reports and Synantec budg reports. A comprehensive list of Axiom group's TTPs obtained through the CRS segment of the threat diagnosis framework is shown in Table 8. This also mustrates that an added contribution of our framework is that it is capable of *data au mentation*. Meaning it can connect threats to TTPs from different data sources in the car of missing connections and provides a comprehensive TTP coverage for each threat jurcident.

TTP	AT &CK	IBM X- Force	Symantec
Zero-day exploits in Internet Explorer 10 to transfer malwar.		$\checkmark$	$\checkmark$
Remote Access Trojan		$\checkmark$	$\checkmark$
Credential Dumping	$\checkmark$		$\checkmark$
Accessibility using RDP	$\checkmark$	$\checkmark$	$\checkmark$
Supply Chain Attacks		$\checkmark$	$\checkmark$
Strategic Web Compromise through waterhol attack		$\checkmark$	$\checkmark$
Deposits malware signed with stolen Bit9 keys		$\checkmark$	$\checkmark$
Payload Obfuscated using XOR encoding	$\checkmark$	$\checkmark$	$\checkmark$
Direct injection of attack code into the International Joren pro- cess without writing it to disk		$\checkmark$	$\checkmark$
The User-Agent header for CnC communication is "lynx"		$\checkmark$	$\checkmark$
CnC traffic as HTTP POST and neH'1." protocol traffic		$\checkmark$	$\checkmark$
Sometimes uses HTTPS (port 44?		$\checkmark$	$\checkmark$
SQL Injection attack on an Internet-1, ing Neb server		$\checkmark$	$\checkmark$
Oracle Java SE Runtime Envir Ament Hotspot code execution		$\checkmark$	$\checkmark$

Table 8: Axiom hacker group TTP Coverage by ATT&CK, 'BM A-Force and Symantec

The ATT&CK taxor omy halo been used as a baseline to model adversary's attack patterns in multiple rescare works. In [68], the ATT&CK taxonomy is used as a baseline to collect historical knowledge of the threat agent for automating the comparison of postcompromise actions of a lversaries. Similarly, in [69], the ATT&CK taxonomy is used to design and implement on automated red teaming system by employing TTPs to drive the atomic actic is of a linelligent decision-making system. To aggregate analysis of the enterprise' log date, [70] uses the ATT&CK taxonomy to model malicious threat scenarios based on adv rsary's attack patterns. It helps select useful audit log data from massive ravelogs. In [71], a visual analytic tool that integrates the analytics of a cyber attack with the data required by the security analyst is proposed. The analytics of cyber attacks are the norm ATT&CK taxonomy which is used to establish an appropriate baseline against which to compare anomalies.

As such, A TT&CK serves as the best candidate to construct our required training dataset. We therefore, extracted all the relevant threats, TTPs and artifacts from the on ... ^TT&CK and constructed a dataset, a small excerpt of which is shown in table 9. The nighlighted data fields represent the presence of the TTP in a threat. The data

	Table 9. An excerpt of constructed ATT&CK MITKE benchma <sup>rb</sup> data.							
TTP	T1001	T1002	T1003	T1004	T1005	Л 'С <sup>0</sup> б		
ID								
$\mathbf{TTP}$	Data Ob-	Data	Credential	Winlogon	Data	ine Sys-	Cyber	
$\mathbf{Title}$	fuscation	Com-	Dumping	Helper	from	+em Logi-	Threat	
		pressed		$\mathrm{DLL}$	Loci '	al Offset	Class	
					System.			
	1	0	1	0	0	0	Axiom	
	1	0	1	0		0	Patchwork	
	1	1	0	0	1	0	Duqu	
	1	0	0	0	0	0	S-Type	
	1	0	1	0	0	0	Oldrea	
	1	1	1	0	v	0	Prikorma	
	0	1	1	0	9	0	Ke3chang	
	0	1	1	0	1	0	APT2	
	0	1	1	0	1	0	TG-3390	
	0	1	0	U	0	0	Lazarus	
	0	1	1	L	0	0	Fin6	

fields are normalized to remove zero entries. A TTD network is then built '... 'ween threat incidents, TTPs and detection mechanisms.

and a farmation of ATTROUGH MITTER

### 7.2. Effectiveness of SIRS

The effectiveness of the SIRS module is evaluated with financial malware TTP dataset and ATT&CK Taxonomy. A TT<sup>+</sup> a. 'ionary for 11 different kinds of malware is created using threat reports of [60], [2<sup>4</sup>, [15] a. d [13]. The TTPs related to different phases of cyber kill chain are given in Table <sup>2</sup>. These TTPs are stored in the TTP dictionary and are queried against the indeled malware and ATT&CK documents. The effectiveness of the SIRS is evaluated by non-asy ring precision and recall for each TTP in the dictionary. Precision measures the malware of correctly retrieved documents shown in equation 6. While recall measures are number of correct items collected by the retrieval system given in equation 7[72].

$$Frecision = \frac{TruePositive}{TruePositive + FalsePositive}$$
(6)

$$iecall = \frac{TruePositive}{TruePositive + FalseNegative}$$
(7)

The av rage of precision and recall for all TTPs is calculated. A weighted harmonic mean of bc h the r arameters, measured using F-measure, is shown in equation 8[72].

$$F - measure = 2 * \frac{Precision * Recall}{Precision + Recall}$$
(8)

The results related to the TTPs of each phase of the cyber kill chain are shown in Tau'e 1'. In can be seen that the overall effectiveness of the SIRS for financial malware documents is quite good with an average F-measure of 0.94.

	Precision		Recall		F-r easure	
TTP Queries (Kill Chain Phases)	POS Malware	ATT&CK	POS Malware	ATT&CK	PC. Malwar	1TT&CK
Reconaissance	95%	22%	100%	33%	0.97	0.26
Weaponize and De- liver	94%	48%	87%	46%	0.9 ,	0.47
Exploit	97%	50%	98%	27%	0.97	0.35
Install	93%	85%	96%	80%	0.8 1	0.82
Command and Control	95%	90%	90%	84%	0.9 2	0.87
Actions and Objec- tives	94%	82%	89%	78° J	0.91	0.79

### 7.3. Effectiveness of TTD Network

To check the sensitivity of methods towards  $a_{+}^{+}a_{+}w_{-}$  evaluated the *TTD Semantic* Network using the hold-out cross validation method [13]. In this method, the given dataset is split into separate parts, i.e., a traning set and a testing set. In our case, we are using the threat-TTP ground truth data provided by ATT&CK MITRE to train our cyber threat prediction model. For the lest set, we examined a set of recent cyber threat incidents listed in ATT&CK MITRE and we determined their threat actors. We searched and downloaded more recent at sets on these threat actors from four different well-known cyber threat sources, i.e., IBM Z-Force [5], Symantec [23], FireEye [24] and CrowdStrike [25]. These new threat non-means thereafter comprised our test set. These threat sources were chosen so that they dnir to some extent in the level of abstraction, the vocabulary used and the level f specialization in the type of threat. Out of these, IBM X-Force supports STIX er coded  $\TIR$ , while for the rest the reports are encoded in STIX before importing them . to the TTP dictionary.

We experimented with so veral magnitude learning algorithms for determining the best prediction accuracy. The esuits are shown in table 11. The results depict that the belief network outperfor is the other techniques which is clearly suggestive that it is more suitable to our roblem. It is easy to see why. When a data breach occurs, security analysts analyze the  $\gamma$  from multiple aspects. For recently occurred incidents the data is likely to nation conflicting or missing feature values as various vendors begin to study it. For c at breaches that have been well studied the features are likely to be well established and complete. Our dataset has a mix of both new and well studied data breaches. Therefore, we required a classifier that can reliably predict a cyber threat f. vil in ne presence of erroneous or missing features. Furthermore, we observed that the L<sup>1</sup> of network is less sensitive to irrelevant features (i.e., noisy data) that an ad ersary an intentionally insert to poison the cyber threat prediction. Finally, a belief ne work c nsiders all the features independently and trains in linear time as compare' to other iterative expensive approximation approaches. We evaluated the accure by of the TTD network for each threat source separately. As a result, we observed a vary g acc racy between the ranges of 53%-100% based on the cyber threat source chosen as the test subject as shown in table 11. The results demonstrate the sensitivity of  $v \in \mathcal{I}$  is semantic network to the amount of TTPs that exist in the test data, which is an sential requirement of our model to achieve accurate classification. It is to note that our results are not a reflection of the data quality of these threat sources as they are based on a very selective CTIR test set and their use of the TTP labes de med in our ATT&CK training dataset. Not all cyber threat sources follow the same  $1.7^{-3}$  vocabulary because of the reasons quoted above. The TTP statistics for cyber threathing in the training and test datasets are shown in figure 13. It can be seen that since and IBM X-force test dataset has the highest number of TTPs, it also has the highest tyber that at prediction accuracy (100%). The Fireeye test dataset has the next highest number of TTPs, thus the prediction accuracy is slightly lower (95%), followed by CrowdStrike and Symantec that are at respectively 67% and 53%. Also it is worth noting that threather threather is a misclassification due to missing TTPs the model picks the threather threather, e.g., in one case "Carbanak", was misclassified as "GCMAN" ooth of which target banks for financial fraud.

|--|

	IBM Force	Х-	Symantec	∿ire'∠ye	$\mathbf{CrowdStrike}$
Belief Network	100%		53%	95%	67%
Decision Tree	30%		18%	33%	33%
Random Forest	45%		35%	57%	33%
Deep Learning	20%		12%	14%	17%
$\mathbf{SVM}$	60%		~%	85%	67%



■ATT&CK ■IBM X-Force ¬Symantec ©FireEye ©CrowdStrike

regure 13: ITP statistics for the under consideration cyber threat sources

Further, we evaluated the accuracy of the *TTD Semantic Network* for two kinds of situations, i.e. lost TTP ratio and spurious TTP ratio. The results are shown in figure 14 and 18. In bot situations, accuracy is compared with the three types of lost and spurious TTP ratio  $\sim$  i.e., random, overlap and disjoint. The random TTP ratio represents an un new loop roach of attack injection encountering both the best case and the worst case of the ent artifact identification. While the overlap and disjoint TTP ratios are biased

based on attack injections representing the worst case and best case of  $f_{\rm ex}$  at artifacts identification respectively. The worst case is the one in which the attackers aim is to misguide the security analyst in threat identification. The overlap in the 1 st TTP ratio represents that the disjoint TTPs of the threat are intentionally removed or disguised to increase the number of false positives in the prediction result. While the disjoint lost TTP ratio represents that the overlapped TTPs of the threat are intentionally removed or disguised to increase the detection accuracy in the presence of a spininum number of detected threat artifacts. Similarly, the overlap in spurious TT ratio represents that the overlap in spurious TT ratio represents that the disjoint threat incidents are intentially ac led to increase the number of false positives in the prediction results. While the disjoint STP ratio represents that the disjoint TTPs of the threat incident  $\varepsilon$  is interval and the disjoint threat incident are intentially ac led to increase the automatic represents that the disjoint TTPs of the threat incident  $\varepsilon$  is interval and the disjoint threat incident  $\varepsilon$  is interval.



Figure 15: Accuracy Results for Spurious TTP Scenario

Une to lot  $\Gamma$ TPs. The average accuracy of a random unbiased situation in case of spurious

TTPs is 95% and in the case of lost TTPs is 87%. For an ideal situation, where all the spurious TTPs are disjoint, the accuracy becomes 100%. While for the worsplace when TTP overlap is high the average accuracy is 92%. Similarly, in the case of lost TTPs an ideal situation where the left over TTPs are disjoint the average accuracy becomes 94%. While for the worst case when the left over TTPs have high overlap, the average accuracy is nearly 76%.

Similarly, the ratio of false positive aka type I errors is measure.' by False Discovery Rate (FDR) for lost TTP and spurious TTP ratio [73]. FDR is defined in equation 9:

$$FDR = \frac{\sum FalsePositive}{\sum PositiveExperimentCatcome}$$
(9)

The results are shown in figure 16 and 17. In both s. Liton, FDR is compared for the three types of lost and spurious TTP ratio, i.e., ra. dom, overlap and disjoint. The results show that the overall FDR of spurious TTPs is low as compared to lost TTPs. For an unbiased random situation, the average FDh for spinious TTPs is 6% while the average FDR for lost TTPs is 19%. For an ideal structure are disjoint, the FDR becomes 0. While for the worst are disjoint, the FDR becomes 0. While for the worst ase when TTP overlap is high the average FDR is 10%. Similarly, in the case of lost TTPs an ideal situation where the left over TTPs are disjoint the average FDR becomes 8%. While for the worst case when the left ones TTPs have high overlap, the average FDR is nearly 30%.



Jigure 16: FDR Results for Lost TTP Scenario

33



Figure 17: FDR Results for Spurious T1. Scenario

### 7.4. Effectiveness of Cyber Threat Prediction

The goal of the proposed threat investigation frame ork is to achieve a highly reliable threat prediction. During the experiments, "here are certain factors observed that affect the prediction's reliability. These factors are the prediction selection probability threshold, the number of detected threat primers (TTPD) and the overlap between TTPs.

#### 7.4.1. Threshold Impact on Prediction Relia ility

Cyber threat investigation results in threat prediction by assigning an appropriate probability value to the threat incident. Among those threat incidents, a prediction set above a certain probability threshort one seeds to be selected. For this purpose, we analyzed the impact of increasing and decreasing probability threshold values over the prediction reliability. We evaluated prediction reliability using the probability thresholds: 25%, 50%, 75% and 100%. In figure 18, an example of one such scenario is given. It can be seen that according to the three to a cordinate the probability, i.e., 10.93% that become 100% of er normalizing the probability using equation 10.

$$x_n = \frac{x - x_{min}}{x_{max} - x_{min}} * 100$$
(10)

If the system h.  $\dot{}$  to select only one threat incident with the highest probability based on the observed TTPs, ben the Decebal malware will be selected. However, in our case, the framework selects a set of threat incidents above a certain probability threshold. In figure 19, the p. c. citic reliability for 10 different threat scenarios is shown. The series represents the threat scenarios. The results show that when we increase the threshold value the prediction reliability increases but there were certain scenarios where no attack classes were identified above 90%. Therefore, there was a need to select a threshold value that alr ays generates a prediction set for all threat scenarios and also provides one with high r liability Based on experimentation, this threshold value is selected as 75%.



Figure 18: An example c'cype \_\_\_\_\_eat prediction



Figure 19: Impact of in asing and decreasing Probability threshold value over Prediction Reliability

### 7.4.2. TTPD mp( et or Prediction Reliability

The number of  $4\epsilon$  ected TTPs, i.e., TTPD also has a great impact on prediction reliability. Durine experiments, it is observed that the prediction reliability is high if there are no redetected TTPs related to the threat. In figure 20, the prediction reliability and TTPD radio are shown. The graph shows that the reliability of the prediction is above 50% if there exists TTPD 50%.



Figure 20: The impact of number of detected artifacts, i. TTPD over reliability

#### 7.4.3. Disjoint and Overlapping TTPs Impact on Prediction Reliability

There are certain TTPs that are more common that, there, e.g., the data obfuscation TTP is more frequently used than routing co. mand and control traffic over a nonstandard custom port bypassing security tools [74]. In the, a significant overlap is usually observed across multiple threats for these common a TPs. We will refer to TTPs that do not occur across different threats in our threat rediction set as disjoint. This particular section details the effect of overlapping as well as disjoint TTPs on the threat prediction reliability. A subset of the dataset is divided into two classes, where one class has threat incidents having mostly disjoint TTPs and the prediction reliability for threat incidents having mostly overlapping TTPs. It is observed that the prediction reliability for threat incidents having disjoint TTPs is about  $20^{\circ\circ}$  higher than the threat incidents having overlapping TTPs. The reason behind the fluct is do joint TTPs depict the distinguishing features of a threat incident that help a threat to be identified quite easily and with a high reliability as compared to the overlapping TTPs. That share less distinguishing features. Figure 21 illustrates the results.



1. July 21: The impact of disjoint and Overlapping TTPs on Prediction Reliability

#### 7.5. Efficiency of Data Breach Investigation Framework

The TTD semantic network is constructed offline or as a background process so we are not interested in the time taken in training. However, to study the . pact of the dataset size on the time for threat detection, we performed experiment. with different dataset sizes. We performed our experiments on a PC with Intel Co end 42100 processors and 1.7 GHz and 2.4 GHz processing speed, 4 GB of RAM rur ing a 64 bit Windows 10. The results can be seen in figure 22. We observed the threat detection time by partitioning the dataset into four equal parts. The observations are recorded for 25%, 50%, 75% and 100% instances of the dataset. Then for each  $\alpha$  taset d vision, the number of observed threat artifacts were incrementally increased. The a. in of observed threat artifacts ranged from 10% to 100%. The results show t' at t' \_ ize of the dataset does not have a significant impact on the threat detection i.e. The factor that matters most is the threat detection time provided the threat "tifa...s are fed into the threat investigation framework. The experimental results in figure 23 for threat prediction time with increasing number of threat artifacts show bat or r system can predict threat incidents for 6 detected TTPs in an average time of 0.04 seconds. The results do not show any significant difference in the detection time when the amber of TTPD increases. The detection time for 34 TTPD is approximately 0.5. conds and the average detection time is 0.15 seconds. Considering that data breach incidents remain undetected for months or even years [15] and attacks such as DDOS 1. t 'or 9-12 hours [75], the running time of our system is quite practical. Moreover, "e ca. claim that the threat detection time of our proposed data breach investigation frame ork is low as compared to existing threat detection mechanisms.







Figure 23: Impact of Dataset size on Threat D. oction Time

### 7.6. Feature Based Comparative Analysis

The performance of the proposed data breach investigation framework cannot be directly compared with any existing threat detection mechanism due to its novelty. However, a comparative feature analysis was conducted. The existing threat detection approaches fall under (1) signature-based,  $(2_1 \text{ nor } a_{2_1})$  (behavior) based and (3) stateful protocol analysis [76]. In Table 12, a feature based detailed comparison of each approach is given.

Features	S <sup>:</sup> ,ла. *e	Anomaly	Stateful Protocol	Proposed Framework
Fixed Signatures	$\sim$	$\checkmark$	$\checkmark$	
Dynamic Signatures		$\checkmark$	$\checkmark$	
Known Attack		$\checkmark$	$\checkmark$	$\checkmark$
Unforeseen Attack		$\checkmark$		$\checkmark$
Semantic Mapping				$\checkmark$
Statistical Prediction		$\checkmark$		$\checkmark$

Table 12: Feature based comparative analysis c. threat diagnosis framework with existing threat detection techniques

The signature-1 ased threat detection approaches work with fixed format signatures that need to be regule. 'v updated to cope with new threats. They can efficiently detect known attacks r rovi'led the same previously known signatures are used for future attacks. As discussed  $\epsilon$ , "lie", the signatures are susceptible to change and have a very short life, thus signature base.' pproaches cannot detect unknown attacks or the same family of attacks using different signatures. On the converse, anomaly-based detection approaches are heurist. or rule based and are capable of detecting new threats by classifying behaviors as  $r_{\rm max}$  abnormal using machine learning and artificial intelligence techniques. They use a fixed patterned. Also, the false platmer ate is high due to weak normal and abnormal behavior profile of the attack trane. The alarm system alerts an abnormal activity without providing further confect an information to guide the security analyst towards further threat investigations. The stateful protocol analysis is the specification based approach that identifies

abnormal sequences of commands in protocol functional trace. This teal ique falsely classifies threats that have the correct protocol execution sequence as been in. The existing threat detection approaches are not capable of mapping the detected network and host artifacts to the adversary's TTPs. The proposed data breach investigation framework works with both semantic and statistical procedures to investigate threat incidents. It does not follow any fixed pattern for threat detection rathed it shap, semantically the detected network and threat artifacts to the adversary's attact, behavior quoted as TTPs in the structured CTIR and unstructured security doed lients. The mapping helps in identifying both known and unforeseen attacks. The framework ither classifies the unforeseen attacks to an existing attack family or suggests fur her investigations to the security analyst to help discover the root cause of the underlying threat activity.

#### 8. Conclusions

In this paper, a novel data breach investigation is mewerk is presented that investigates cyber threat incidents using high-level adverse, v's 1 l'Ps from cyber threat intelligence documents. At the core of the framework line the TTD semantic network that is based on the idea of semantically mapping low-ic al threat artifacts to high-level adversarial attack techniques and employs them to reason about an incident occurrence. The methodology and results presented highlight that security incidents can be mapped to tactics, that are further mapped to  $\operatorname{arti}^{f}$  at a way that machines can identify these connections with certain probabilities. How for this research paves the way for cyber security investigations with partial or in amplete information. The system extracts TTPs from CTIR and employs them to Grid semantic relevance (using LSI) between the threats and the detection mechanisms. The TTD semantic network is enriched by training a belief network that maps the care to TTPs to predict the most probable set of threat incidents based on the detected threat a rtifacts provided by the security analyst. In the case of an unforeseen threat, the s<sub>c</sub> ter helps the security analyst investigate the threat artifacts against the most p oba' le attack family by suggesting the most optimized and cost-effective detection me. bar .sms This mapping of the threat artifacts to the TTPs was able to detect the ' neat ... dents with high accuracy and low false positives in the case of lost and spin ous detected threat artifacts. The TTD semantic network is extensible as more threat doc, ments arrive they can be indexed and ranked using LSI. Similarly, the detec<sup>+</sup> on nechanisms can also be dynamically updated as the ATT&CK taxonomy updates TF e evaluation results demonstrate that the detection time of the framework is qui'e low compared to the considerable time it typically takes to investigate data breaching dents. In the future, we will integrate and automate the mitigations for the identifies + irea' incidents.

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# **Highlights**

- A novel framework for cyber data breach investigation.
- Predicts cyber-threat with high accuracy.
- Employs high-level Tactics Techniques and Procedures for cyber u. .t investigation.
- A first of its kind framework that uses a comprehensive vocabulary for structured TTP analysis.